

## THE DISTRIBUTION OF VALUES IN THE QUADRATIC ASSIGNMENT PROBLEM

ALEXANDER BARVINOK AND TAMON STEPHEN

We obtain a number of results regarding the distribution of values of a quadratic function  $f$  on the set of  $n \times n$  permutation matrices (identified with the symmetric group  $S_n$ ) around its optimum (minimum or maximum). We estimate the fraction of permutations  $\sigma$  such that  $f(\sigma)$  lies within a given neighborhood of the optimal value of  $f$  and relate the optimal value with the average value of  $f$  over a neighborhood of the optimal permutation. We describe a natural class of functions (which includes, for example, the objective function in the Traveling Salesman Problem) with a relative abundance of near-optimal permutations. Also, we identify a large class of functions  $f$  with the property that permutations close to the optimal permutation in the Hamming metric of  $S_n$  tend to produce near optimal values of  $f$  (such is, for example, the objective function in the symmetric Traveling Salesman Problem). We show that for general  $f$ , just the opposite behavior may take place: an average permutation in the vicinity of the optimal permutation may be much worse than an average permutation in the whole group  $S_n$ .

**1. Introduction.** The Quadratic Assignment Problem (QAP for short) is an optimization problem on the symmetric group  $S_n$  of  $n!$  permutations of an  $n$ -element set. The QAP is one of the hardest problems of combinatorial optimization, whose special cases include the Traveling Salesman Problem (TSP) among other interesting problems.

Recently, the QAP has been of interest to many people. An excellent survey is found in Burkard et al. (1999). Despite this work, it is still extremely difficult to solve QAP's of size  $n = 20$  to optimality, and the solution to a QAP of size  $n = 30$  is considered noteworthy (see, for example, Anstreicher et al. 2002 and Brügger et al. 1998). There are some approximability results known for special classes of QAPs (Arkin et al. 2001), but mostly "bad news" (nonapproximability results) for the general case (Ausiello et al. 1999).

The goal of this paper is to study the distribution of values of the objective function of the QAP. We hope that our results would allow one, on one hand, to understand the behavior of various heuristics, and, on the other hand, to estimate the optimum using some simple algorithms based on random or partial enumeration with guaranteed complexity bounds. In particular, we estimate how well the sample optimum from a random sample of a given size approximates the global optimum.

**1.1. The Quadratic Assignment Problem.** Let  $\text{Mat}_n$  be the vector space of all real  $n \times n$  matrices  $A = (a_{ij})$ ,  $1 \leq i, j \leq n$  and let  $S_n$  be the group of all permutations  $\sigma$  of the set  $\{1, \dots, n\}$ . There is an action of  $S_n$  on the space  $\text{Mat}_n$  by simultaneous permutations of rows and columns: we let  $\sigma(A) = B$ , where  $A = (a_{ij})$  and  $B = (b_{ij})$ , provided  $b_{\sigma(i)\sigma(j)} = a_{ij}$  for all  $i, j = 1, \dots, n$ . One can check that  $(\sigma\tau)A = \sigma(\tau A)$  for any two permutations  $\sigma$  and  $\tau$ . There is a standard scalar product on  $\text{Mat}_n$ :

$$\langle A, B \rangle = \sum_{i,j=1}^n a_{ij}b_{ij} \quad \text{where } A = (a_{ij}) \text{ and } B = (b_{ij}).$$

Received July 17, 2001; revised October 9, 2002.

*MSC 2000 subject classification.* Primary: 90C27, 90C59, 68R05, 05E10.

*ORMS subject classification.* Primary: Analysis of algorithms; secondary: Networks/graphs, traveling salesman.

*Key words.* Quadratic Assignment Problem, distribution, symmetric group, randomized algorithms, heuristics, representation theory.

Let us fix two matrices  $A = (a_{ij})$  and  $B = (b_{ij})$  and let us consider a real-valued function  $f: S_n \rightarrow \mathbb{R}$  defined by

$$(1.1.1) \quad f(\sigma) = \langle B, \sigma(A) \rangle = \sum_{i,j=1}^n b_{\sigma(i)\sigma(j)} a_{ij} = \sum_{i,j=1}^n b_{ij} a_{\sigma^{-1}(i)\sigma^{-1}(j)}.$$

The problem of finding a permutation  $\sigma$  where the maximum or minimum value of  $f$  is attained is known as the *Quadratic Assignment Problem* (QAP). It is one of the hardest problems of combinatorial optimization. From now on we assume that  $n \geq 4$ .

We contrast the QAP with the *Linear Assignment Problem* (LAP) of maximizing

$$f(\sigma) = \sum_{i=1}^n d_{i\sigma(i)}$$

for a given  $n \times n$  matrix  $D = (d_{ij})$ . The LAP is again an optimization problem on the set of permutations. However, there is a well-known polynomial time algorithm for the LAP described, for example, in Papadimitriou and Steiglitz (1982).

Our approach produces essentially identical results for a more general problem.

**1.2. The generalized problem.** Suppose we are given a 4-dimensional array (tensor)  $C = \{c_{kl}^{ij} : 1 \leq i, j, k, l \leq n\}$  of  $n^4$  real numbers and the function  $f$  is defined by

$$(1.2.1) \quad f(\sigma) = \sum_{i,j=1}^n c_{\sigma(i)\sigma(j)}^{ij}.$$

If  $c_{kl}^{ij} = a_{ij}b_{kl}$  for some matrices  $A = (a_{ij})$  and  $B = (b_{kl})$ , we get the special case (1.1.1) we started with. The convenience of working with the generalized problem is that the set of objective functions (1.2.1) is a vector space. This problem was considered in Lawler (1963).

We introduce the standard Hamming metric on the symmetric group  $S_n$ .

1.3. DEFINITION. For two permutations  $\sigma, \tau \in S_n$ , let the distance  $\text{dist}(\sigma, \tau)$  be the number of indices  $1 \leq i \leq n$  where  $\sigma$  and  $\tau$  disagree:

$$\text{dist}(\sigma, \tau) = |i : \sigma(i) \neq \tau(i)|.$$

One can observe that the distance is invariant under the left and right actions of  $S_n$ :

$$\text{dist}(\sigma\sigma_1, \sigma\sigma_2) = \text{dist}(\sigma_1, \sigma_2) = \text{dist}(\sigma_1\sigma, \sigma_2\sigma)$$

for all  $\sigma_1, \sigma_2, \sigma \in S_n$ .

For a permutation  $\tau$  and an integer  $k \geq 0$ , we consider the “ $k$ th sphere” around  $\tau$ :

$$U(\tau, k) = \{\sigma \in S_n : \text{dist}(\sigma, \tau) = n - k\}.$$

Hence, for any permutation  $\tau$  the group  $S_n$  splits into the disjoint union of  $n$  spheres  $U(\tau, k)$  for  $k = 0, 1, \dots, n-2, n$ . The sphere  $U(\tau, k)$  consists of the permutations that agree with  $\tau$  on precisely  $k$  symbols.

Let  $f: S_n \rightarrow \mathbb{R}$  be a function of type (1.1.1) or (1.2.1). Let

$$\bar{f} = \frac{1}{n!} \sum_{\sigma \in S_n} f(\sigma)$$

be the average value of  $f$  on the symmetric group and let

$$f_0 = f - \bar{f}$$

be the “shifted” function. Hence, the average value of  $f_0$  is 0. Let  $\tau$  be a permutation where the maximum value of  $f_0$  is attained, so  $f_0(\tau) \geq f_0(\sigma)$  for all  $\sigma \in S_n$  and  $f_0(\tau) > 0$  unless  $f_0 \equiv 0$ . Because we measure the deviation from 0, we can analyze minimization problems by replacing  $f_0$  with  $-f_0$  and studying the related maximization problem. We maintain the definitions of  $\bar{f}$ ,  $f_0$ , and  $\tau$  in future sections.

We are interested in the following questions:

- Given a constant  $0 < \gamma < 1$ , how many permutations  $\sigma \in S_n$  satisfy  $f_0(\sigma) \geq \gamma f_0(\tau)$ ? In particular, how well does the sample optimum of a set of randomly chosen permutations approximate the true optimum?

- How does the average value of  $f_0$  over the  $k$ th sphere  $U(\tau, k)$  compare with the optimal value  $f_0(\tau)$ ? In particular, is a random permutation from the vicinity of the optimal permutation better than a random permutation from the whole group  $S_n$ ?

We remark that it is easy to compute the average value  $\bar{f}$  (see Lemma 6.1).

This paper is organized as follows. In §§2–5 we state our main results for the general QAP (§5) and its particular cases (§§2–4), and in §§6–11 we provide proofs. In §12, we formulate some open problems.

1.4. NOTATION. For an integer  $m \geq 0$ , let

$$d_m = \sum_{j=0}^m (-1)^j \frac{1}{j!}.$$

For  $m \geq 2$ , let

$$\nu(m) = \frac{d_{m-2}}{d_m}.$$

It is convenient to agree that  $\nu(0) = 0$  (and  $\nu(1)$  is not defined). One can see that

$$\lim_{m \rightarrow +\infty} \nu(m) = 1$$

and that  $0 \leq \nu(m) < 1$  if  $m$  is odd and  $\nu(m) > 1$  if  $m \geq 2$  is even.

The following functions  $p$  and  $t$  on the symmetric group  $S_n$  play a special role in our approach. For a permutation  $\sigma \in S_n$ , let

$$p(\sigma) = |i : \sigma(i) = i|$$

be the number of fixed points of  $\sigma$  and let

$$t(\sigma) = |i < j : \sigma(i) = j \text{ and } \sigma(j) = i|$$

be the number of 2-cycles in  $\sigma$ .

One can show that  $p(\sigma)$ ,  $p^2(\sigma)$ , and  $t(\sigma)$  are functions of type (1.2.1) for some particular tensors  $\{c_{kl}^{ij}\}$ , (see Remark 7.6). We denote by  $\varepsilon$  the identity permutation in  $S_n$  and by  $|X|$  the cardinality of a set  $X$ .

**2. The bullseye case.** Suppose that the matrix  $A = (a_{ij})$  in (1.1) is symmetric and has constant row and column sums and a constant diagonal

$$\begin{aligned} & a_{ij} = a_{ji} \quad \text{for all } 1 \leq i, j \leq n, \\ \text{for some } a, & \quad \left\{ \begin{array}{l} \sum_{i=1}^n a_{ij} = a \quad \text{for all } j = 1, \dots, n \quad \text{and} \\ \sum_{j=1}^n a_{ij} = a \quad \text{for all } i = 1, \dots, n, \end{array} \right. \\ & a_{ii} = b \quad \text{for some } b \text{ and all } i = 1, \dots, n. \end{aligned}$$

For example,

$$A = \begin{pmatrix} 0 & 1 & 0 & \dots & 0 & 1 \\ 1 & 0 & 1 & \dots & 0 & 0 \\ 0 & 1 & 0 & \dots & 0 & 0 \\ \dots & \dots & \dots & \dots & \dots & \dots \\ 0 & 0 & 0 & \dots & 0 & 1 \\ 1 & 0 & 0 & \dots & 1 & 0 \end{pmatrix}, \quad a_{ij} = \begin{cases} 1 & \text{if } |i-j| = 1 \pmod n, \\ 0 & \text{otherwise,} \end{cases}$$

satisfies these properties and the corresponding optimization problem is the *Symmetric Traveling Salesman Problem*.

Similarly, for the generalized problem (1.2) we assume that for any  $k$  and  $l$  the matrix  $A = (a_{ij})$ , where  $a_{ij} = c_{kl}^{ij}$ , is symmetric with constant row and column sums and has a constant diagonal.

It turns out that the optimum has a characteristic “bullseye” feature with respect to the averages over the sphere  $U(\tau, k)$  (see Definition 1.3).

2.1. THEOREM. *Let*

$$\alpha(n, k) = \frac{k^2 - 3k + \nu(n - k)}{n^2 - 3n},$$

where  $k = 0, 1, \dots, n - 2, n$  and  $\nu$  is the function defined in (1.4). Then, we have

$$\frac{1}{|U(\tau, k)|} \sum_{\sigma \in U(\tau, k)} f_0(\sigma) = \alpha(n, k) f_0(\tau)$$

for  $k = 0, 1, \dots, n - 2, n$ .

In fact, this result holds even if we do not require  $\tau$  to maximize  $f_0$ .

**2.2. The “bullseye” distribution.** We observe that as the sphere  $U(\tau, k)$  contracts to the optimal permutation  $\tau$  (hence  $k$  increases), the *average* value of  $f$  on the sphere steadily improves (as long as  $k \geq 3$ ); see Figure 1.

It is easy to construct examples where *some* values of  $f$  in a very small neighborhood of the optimum are particularly bad, but as follows from Theorem 2.1, such values are relatively rare. In our opinion, this suggests that this special case may be more amenable to local search type heuristics (see, for example, Burkard et al. 1999) than the general

Distribution of values of the objective function with respect to the Hamming distance from the maximum point

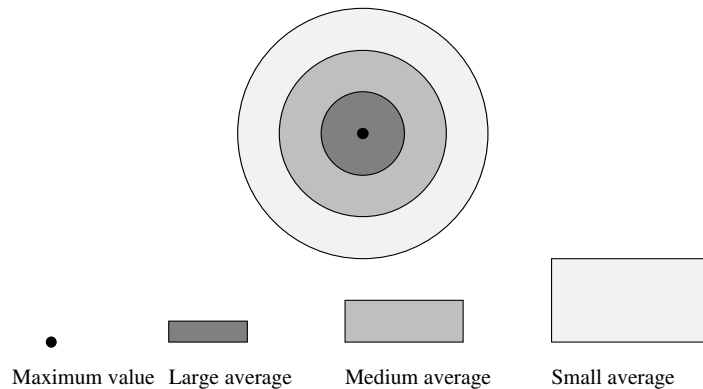


FIGURE 1. Bullseye.

symmetric case (cf. §4). Incidentally, one can observe the same type of the “bullseye” behavior for the LAP and some other polynomially solvable problems, such as the weighted Matching Problem.

Estimating the size of the sphere  $U(\tau, k)$ , we get the following result.

2.3. THEOREM. *Let us choose an integer  $3 \leq k \leq n-3$  and a number  $0 < \gamma < 1$  and let*

$$\beta(n, k) = \frac{k^2 - 3k}{n^2 - 3n}.$$

*The probability that a random permutation  $\sigma \in S_n$  satisfies the inequality*

$$f_0(\sigma) \geq \gamma\beta(n, k)f_0(\tau)$$

*is at least*

$$\frac{(1 - \gamma)\beta(n, k)}{3k!}.$$

We prove our results in §8.

**3. The pure case.** In this section, we consider a more general case of a not necessarily symmetric matrix  $A$  in (1.1) having constant row and column sums and a constant diagonal:

$$\begin{aligned} \text{for some } a \quad & \begin{cases} \sum_{i=1}^n a_{ij} = a & \text{for all } j = 1, \dots, n \text{ and} \\ \sum_{j=1}^n a_{ij} = a & \text{for all } i = 1, \dots, n, \end{cases} \\ a_{ii} = b & \quad \text{for some } b \text{ and all } i = 1, \dots, n. \end{aligned}$$

For example, matrix

$$A = \begin{pmatrix} 0 & 1 & 0 & \dots & \dots & \dots & 0 \\ 0 & 0 & 1 & 0 & \dots & \dots & 0 \\ \dots & \dots & \dots & \dots & \dots & \dots & \dots \\ 0 & 0 & \dots & \dots & \dots & 0 & 1 \\ 1 & 0 & \dots & \dots & \dots & \dots & 0 \end{pmatrix}, \quad a_{ij} = \begin{cases} 1 & \text{if } j = i + 1 \pmod n, \\ 0 & \text{otherwise,} \end{cases}$$

satisfies these properties and the corresponding optimization problem is the *Asymmetric Traveling Salesman Problem*.

Similarly, for generalized problems (see 1.2) we assume that for any  $k$  and  $l$  the matrix  $A = (a_{ij})$ , where  $a_{ij} = c_{kl}^{ij}$  has constant row and column sums and has a constant diagonal. These conditions can be generalized further (see Stephen 2002).

We call this case pure, because as we remark in §§7 and 9, the objective function  $f$  lacks the component attributed to the LAP. More generally, an arbitrary objective function  $f$  in the QAP can be represented as a sum  $f = f_1 + f_2$ , where  $f_1$  is the objective function in a LAP and  $f_2$  is the objective function in some pure case.

The behavior of averages of  $f_0$  over the sphere  $U(\tau, k)$  is described by the following result.

3.1. THEOREM. *Let us define three functions of  $n$  and  $k$ :*

$$\alpha_1(n, k) = 1 - \nu(n - k),$$

$$\alpha_{2e}(n, k) = \frac{k^2 - 3k - n - 3\nu(n - k) + \nu(n - k)n + 4}{n^2 - 4n + 4}, \quad \text{and}$$

$$\alpha_{2o}(n, k) = \frac{k^2 - 3k - n - 2\nu(n - k) + \nu(n - k)n + 3}{n^2 - 4n + 3},$$

where  $k = 0, 1, \dots, n - 2, n$  and  $\nu$  is the function defined in (1.4).

If  $n$  is even, then for some nonnegative  $\gamma_1$  and  $\gamma_2$  such that  $\gamma_1 + \gamma_2 = 1$ , we have

$$\frac{1}{|U(\tau, k)|} \sum_{\sigma \in U(\tau, k)} f_0(\sigma) = (\gamma_1 \alpha_1(n, k) + \gamma_2 \alpha_{2e}(n, k)) f_0(\tau)$$

for  $k = 0, 1, \dots, n - 2, n$ .

If  $n$  is odd, then for some nonnegative  $\gamma_1$  and  $\gamma_2$  such that  $\gamma_1 + \gamma_2 = 1$ , we have

$$\frac{1}{|U(\tau, k)|} \sum_{\sigma \in U(\tau, k)} f_0(\sigma) = (\gamma_1 \alpha_1(n, k) + \gamma_2 \alpha_{2o}(n, k)) f_0(\tau)$$

for  $k = 0, 1, \dots, n - 2, n$ .

Unlike Theorem 2.1, Theorem 3.1 may not hold if  $\tau$  is not the optimal permutation.

It follows from our proof (see Remark 9.3) that for any choice of  $\gamma_1, \gamma_2 \geq 0$  such that  $\gamma_1 + \gamma_2 = 1$ , there is a function  $f$  of type (1.2.1) for which the averages of  $f_0$  over  $U(\tau, k)$  are given by the formulas of Theorem 3.1.

We observe that there are two extreme cases. If  $\gamma_1 = 0$  and  $\gamma_2 = 1$ , then  $f$  exhibits a bullseye type distribution of §2.2. If  $\gamma_1 = 1$  and  $\gamma_2 = 0$ , then  $f$  exhibits a “damped oscillator” type of distribution: the average value of  $f_0$  over  $U(\tau, k)$  changes its sign with the parity of  $k$  and approaches 0 fast as  $k$  gets smaller. In short, if  $f$  has a damped oscillator distribution, there is no particular advantage in choosing a permutation in the vicinity of the optimal permutation  $\tau$ ; see Figure 2.

For a typical function  $f$  one can expect both  $\gamma_1$  and  $\gamma_2$  positive, so  $f$  would show a “diluted” bullseye distribution: the average value of  $f_0$  over  $U(\tau, k)$  improves moderately as  $k$  increases, but not as dramatically as in the bullseye case of §2.

Still, it turns out that we can find sufficiently many reasonably good permutations in the vicinity of the optimal permutation.

Distribution of values of the objective function with respect to the Hamming distance from the maximum point

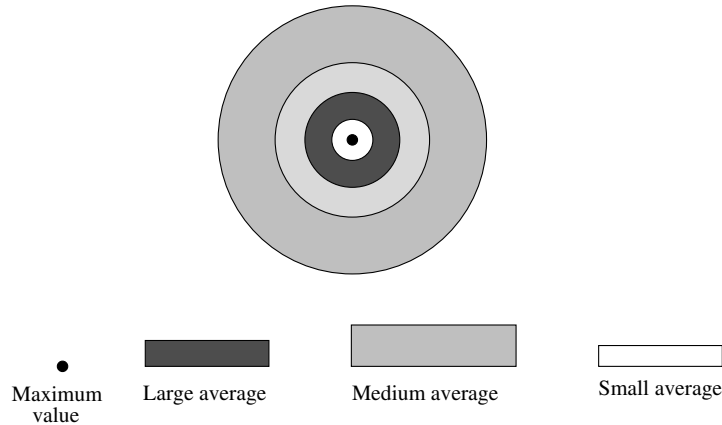


FIGURE 2. Damped oscillator.

3.2. THEOREM. *Let us choose an integer  $3 \leq k \leq n-3$  and a number  $0 < \gamma < 1$ , and let*

$$\beta(n, k) = \frac{k^2 - 3k + 1}{n^2 - 3n + 1}.$$

*The probability that a random permutation  $\sigma \in S_n$  satisfies the inequality*

$$f_0(\sigma) \geq \gamma\beta(n, k)f_0(\tau)$$

*is at least*

$$\frac{(1 - \gamma)\beta(n, k)}{10k!}.$$

Moreover, the statement of the theorem holds if we require, additionally, that  $\sigma \in U(\tau, k)$  (see Remark 9.5).

In particular, by choosing an appropriate  $k$ , we obtain the following corollary.

3.3. COROLLARY. (1) *Let us fix any  $\alpha > 1$ . Then, there exists  $\delta = \delta(\alpha) > 0$  such that for all sufficiently large  $n \geq N(\alpha)$  the probability that a random permutation  $\sigma$  in  $S_n$  satisfies the inequality*

$$f_0(\sigma) \geq \frac{\alpha}{n^2}f_0(\tau)$$

*is at least  $\delta n^{-2}$ . In particular, one can choose  $\delta = \exp\{-c\sqrt{\alpha} \ln \alpha\}$  for some absolute constant  $c > 0$ .*

(2) *Let us fix any  $\epsilon > 0$ . Then, there exists  $\delta = \delta(\epsilon) < 1$  such that for all sufficiently large  $n \geq N(\epsilon)$  the probability that a random permutation  $\sigma$  in  $S_n$  satisfies the inequality*

$$f_0(\sigma) \geq n^{-\epsilon}f_0(\tau)$$

*is at least  $\exp\{-n^\delta\}$ . In particular, one can choose any  $\delta > 1 - \epsilon/2$ .  $\square$*

It follows from Corollary 3.3 that to get a permutation  $\sigma$  which satisfies (1) for any fixed  $\alpha$ , we can use the following straightforward randomized algorithm: sample  $\omega(n)n^2$  random permutations  $\sigma \in S_n$ , where  $\omega(n) \rightarrow +\infty$  arbitrarily slowly, compute the value of  $f$ , and choose the best permutation. With probability tending to 1 as  $n \rightarrow +\infty$ , we will hit the right permutation. If we are willing to settle for a mildly exponential sized sample of the type  $\exp\{n^\beta\}$  for some  $\beta < 1$ , we can achieve a better approximation (2) by searching through a set of randomly selected  $\exp\{n^\beta\}$  permutations. We remark that no algorithm solving the QAP (even in the special case considered in this section) with an exponential in  $n$  complexity  $\exp\{O(n)\}$  is known, although there is a dynamic programming algorithm solving the TSP in  $\exp\{O(n)\}$  time.

We prove our results in §9.

**4. The symmetric case.** In this section, we assume that the matrix  $A = (a_{ij})$ . Subsection 1.1 is symmetric, that is

$$a_{ij} = a_{ji} \quad \text{for all } 1 \leq i, j \leq n.$$

Similarly, in the generalized problem (1.2) we assume that for any  $k$  and  $l$  the matrix  $A = (a_{ij})$ , where  $a_{ij} = c_{kl}^{ij}$ , is symmetric.

Overall, the distribution of values of  $f$  turns out to be much more complicated than in the special cases described in §§2 and 3.

4.1. THEOREM. *Let us define three functions of  $n$  and  $k$ :*

$$\alpha_1(n, k) = \frac{2nk - 2n - k^2 - 3k - \nu(n - k) + 6}{n^2 - 5n + 6},$$

$$\alpha_{2e}(n, k) = \frac{-nk + n + k^2 + k + \nu(n - k) - 4}{2n - 4}, \quad \text{and}$$

$$\alpha_{2o}(n, k) = \frac{-n^2k + nk^2 + n^2 + nk + n\nu(n - k) - 4n - 3k + 3}{2n^2 - 7n + 3},$$

where  $k = 0, 1, \dots, n - 2, n$  and  $\nu$  is the function defined in (1.4).

If  $n$  is even, then for some nonnegative  $\gamma_1$  and  $\gamma_2$  such that  $\gamma_1 + \gamma_2 = 1$ , we have

$$\frac{1}{|U(\tau, k)|} \sum_{\sigma \in U(\tau, k)} f_0(\sigma) = (\gamma_1 \alpha_1(n, k) + \gamma_2 \alpha_{2e}(n, k)) f_0(\tau)$$

for  $k = 0, 1, \dots, n - 2, n$ .

If  $n$  is odd, then for some nonnegative  $\gamma_1$  and  $\gamma_2$  such that  $\gamma_1 + \gamma_2 = 1$ , we have

$$\frac{1}{|U(\tau, k)|} \sum_{\sigma \in U(\tau, k)} f_0(\sigma) = (\gamma_1 \alpha_1(n, k) + \gamma_2 \alpha_{2o}(n, k)) f_0(\tau)$$

for  $k = 0, 1, \dots, n - 2, n$ .

It follows from our proof (see Remark 10.3) that for any choice of  $\gamma_1, \gamma_2 \geq 0$  such that  $\gamma_1 + \gamma_2 = 1$ , there is a function  $f$  of type (1.2.1) for which the averages of  $f_0$  over  $U(\tau, k)$  are given by the formulas of Theorem 4.1. Moreover, at least for even  $n$ , one can choose  $f$  to be a function of type (1.1.1), but we don't prove that here (it is included in Stephen 2002).

**4.2. The “spike” distribution.** As in §3, we see that there are two extreme cases. If  $\gamma_1 = 1$  and  $\gamma_2 = 0$ , then  $f$  has a bullseye-type distribution described in §2.2. If  $\gamma_1 = 0$  and  $\gamma_2 = 1$ , then  $f$  has what we call a “spike” distribution; see Figure 3.

In this case, for  $2 \leq k \leq n - 3$  the average value of  $f_0$  over  $U(\tau, k)$  is negative. Thus, an average permutation  $\sigma \in U(\tau, n - 3)$  presents us with a worse choice than the average permutation in  $S_n$ . However, the average value of  $f_0$  over  $U(\tau, 0)$  is about one half of the maximum value  $f_0(\tau)$ . Thus, there are plenty of reasonably good permutations very far from  $\tau$  and we can easily get such a permutation by random sampling.

Distribution of values of the objective function with respect to the Hamming distance from the maximum point

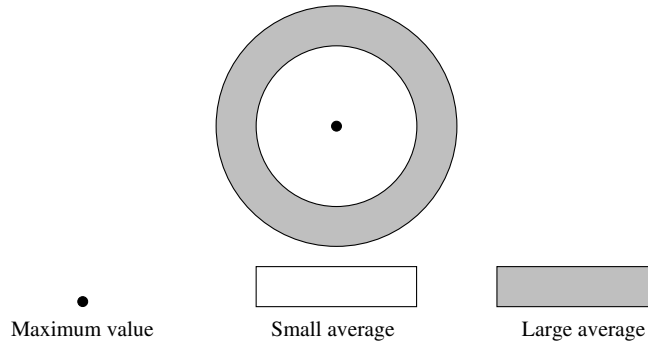


FIGURE 3. Spike.



The distribution of a typical function  $f$  around its optimal permutation in the symmetric QAP is a certain mixture of the bullseye and spike distributions. This interference of the bullseye and spike distributions (which, in some sense, are “pulling in the opposite directions”) provides, in our opinion, a plausible explanation of the computational hardness of the general symmetric QAP even in comparison with other NP-hard problems such as the TSP.

We obtain the following estimate for the number of near-optimal permutations.

4.3. THEOREM. *Let us choose an integer  $3 \leq k \leq n-3$  and a number  $0 < \gamma < 1$ , and let*

$$\beta(n, k) = \frac{3k-5}{n^2 - kn + k + 2n - 5}.$$

*The probability that a random permutation  $\sigma \in S_n$  satisfies the inequality*

$$f_0(\sigma) \geq \gamma\beta(n, k)f_0(\tau)$$

*is at least*

$$\frac{(1-\gamma)\beta(n, k)}{5k!2^k}.$$

One can notice that the obtained bound is essentially weaker than the bounds of Theorems 2.3 and 3.2. In Stephen (2002), we show that at least for the generalized problem (1.2), the bounds of §§2 and 3, and Corollary 3.3, in particular, do not extend. The question of whether the estimates can be improved for problem (1.1) remains open. We prove the results of this section in §10.

**5. The general case.** It appears that the difference between the general case of problems (1.1) and (1.2) and the symmetric case of §4 is not as substantial as the difference between the symmetric case and the special cases of §§2 and 3.

First, we describe the behavior of averages of  $f_0$  over the spheres  $U(\tau, k)$ .

5.1. THEOREM. *Let us define five functions of  $n$  and  $k$ :*

$$\begin{aligned} \alpha_1(n, k) &= \frac{-k(n-k) + n - 2}{n - 2}, \\ \alpha_2(n, k) &= 1 - \nu(n-k), \\ \alpha_3(n, k) &= \frac{2nk - 3k - 2n - k^2 - \nu(n-k) + 6}{n^2 - 5n + 6}, \\ \alpha_4(n, k) &= \frac{k + \nu(n-k) - 2}{n - 2}, \quad \text{and} \\ \alpha_{5_0}(n, k) &= \frac{-2nk + 3k^2 - 3k + \nu(n-k)n + n - 3}{n^2 - 2n - 3}, \end{aligned}$$

where  $k = 0, 1, \dots, n-2, n$  and  $\nu$  is the function of Subsection 1.4.

If  $n$  is even, then for some nonnegative  $\gamma_1, \gamma_2, \gamma_3$ , and  $\gamma_4$  such that  $\gamma_1 + \gamma_2 + \gamma_3 + \gamma_4 = 1$ , we have

$$\frac{1}{|U(\tau, k)|} \sum_{\sigma \in U(\tau, k)} f_0(\sigma) = (\gamma_1\alpha_1(n, k) + \gamma_2\alpha_2(n, k) + \gamma_3\alpha_3(n, k) + \gamma_4\alpha_4(n, k)) f_0(\tau)$$

for  $k = 0, 1, \dots, n-2, n$ .

If  $n$  is odd, then for some nonnegative  $\gamma_1, \gamma_2, \gamma_3, \gamma_4,$  and  $\gamma_5$  such that  $\gamma_1 + \gamma_2 + \gamma_3 + \gamma_4 + \gamma_5 = 1,$  we have

$$\frac{1}{|U(\tau, k)|} \sum_{\sigma \in U(\tau, k)} f_0(\sigma) = (\gamma_1 \alpha_1(n, k) + \gamma_2 \alpha_2(n, k) + \gamma_3 \alpha_3(n, k) + \gamma_4 \alpha_4(n, k) + \gamma_5 \alpha_{5_o}(n, k)) f_0(\tau)$$

for  $k = 0, 1, \dots, n - 2, n.$

It follows from our proof (see Remark 11.3) that for any choice of  $\gamma_1, \dots, \gamma_4 \geq 0$  ( $n$  even) or  $\gamma_1, \dots, \gamma_5 \geq 0$  ( $n$  odd) summing up to 1, there is a function  $f$  of type (1.2.1) for which the averages of  $f_0$  over  $U(\tau, k)$  are given by the formulas of Theorem 5.1. Moreover, at least for even  $n,$  one can choose  $f$  to be a function of type (1.1.1), but we don't prove that here (see Stephen 2002).

Letting all but one  $\gamma$  equal to 0, we obtain various extreme distributions: the bullseye (when  $\gamma_3 = 1$  or  $\gamma_4 = 1,$  cf. §2), damped oscillator (when  $\gamma_2 = 1,$  cf. §3), and spike (when  $\gamma_1 = 1$  or  $\gamma_{5_o} = 1,$  cf. §4) types. We do not get any new type of distribution, but we can find a sharper spike than in the symmetric case.

**5.2. Example of the sharp spike.** Let us consider the function

$$f(\sigma) = \frac{-p(\sigma)(n - p(\sigma)) + n - 2}{n - 2},$$

where  $p(\sigma)$  is the number of fixed points of  $\sigma$  (see §1.4). In §11 (cf. Remark 11.3), we show that  $f$  is a function of type (1.2.1) and that  $\bar{f} = 0.$  The maximum value of 1 is attained at the identity and at the permutations without fixed points. All other values of  $f$  are negative; see Figure 4.

One can construct a function of type (1.1.1) whose average values over the spheres  $U(\varepsilon, k),$  where  $\varepsilon$  is the identity permutation, coincide with those for  $f.$

Our bound for the number of nearly optimal permutations is only slightly weaker than the bound of Theorem 4.3 in the symmetric case.

**5.3. THEOREM.** *Let us choose an integer  $3 \leq k \leq n - 3$  and a number  $0 < \gamma < 1,$  and let*

$$\beta(n, k) = \frac{k - 2}{n^2 - nk + k - 2}.$$

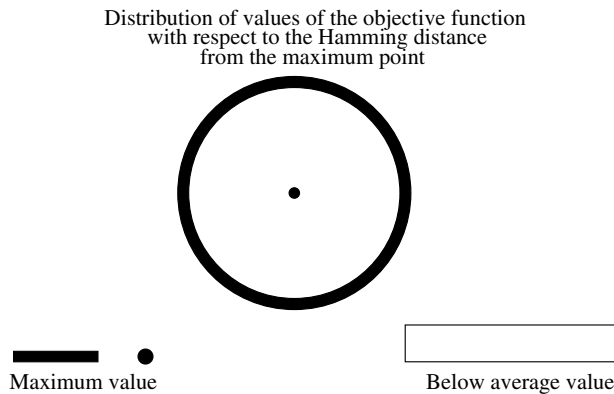


FIGURE 4. Sharp spike.

The probability that a random permutation  $\sigma \in S_n$  satisfies

$$f_0(\sigma) \geq \gamma\beta(k, n)f_0(\tau)$$

is at least

$$\frac{(1-\gamma)\beta(k, n)}{5k!}.$$

By choosing an appropriate  $k$ , we obtain the following corollary.

**5.4. COROLLARY.** (1) *Let us fix any  $\alpha > 1$ . Then there exists  $\delta = \delta(\alpha) > 0$  such that for all sufficiently large  $n \geq N(\alpha)$ , the probability that a random permutation  $\sigma$  in  $S_n$  satisfies the inequality*

$$f_0(\sigma) \geq \frac{\alpha}{n^2}f_0(\tau)$$

is at least  $\delta n^{-2}$ . In particular, one can choose  $\delta = \exp\{-c\alpha \ln \alpha\}$  for some absolute constant  $c > 0$ .

(2) *Let us fix any  $\epsilon > 0$ . Then there exists  $\delta = \delta(\epsilon) < 1$  such that for all sufficiently large  $n \geq N(\epsilon)$ , the probability that a random permutation  $\sigma$  in  $S_n$  satisfies the inequality*

$$f_0(\sigma) \geq n^{-1-\epsilon}f_0(\tau)$$

is at least  $\exp\{-n^\delta\}$ . In particular, one can choose any  $\delta > 1 - \epsilon$ .  $\square$

As in §3, we conclude that for any fixed  $\alpha > 1$ , the randomized algorithm of sampling  $\omega(n)n^2$  permutations, where  $\omega(n) \rightarrow +\infty$  arbitrarily slowly, produces a permutation  $\sigma$  satisfying (1) with probability tending to 1 as  $n \rightarrow +\infty$ . This simple randomized algorithm produces a better bound in a more general situation than some known deterministic algorithms based on semidefinite relaxations of the QAP (cf. Ye 1999 and references therein).

If we are willing to settle for an algorithm of mildly exponential complexity, we can achieve the bound of type (2), which is weaker than the corresponding bound of Corollary 3.3.

We prove our results in §11.

**6. Preliminaries.** First, we show that it is indeed easy to compute the average value  $\bar{f}$  of a function  $f$  defined by (1.1.1) or (1.2.1). The result is not new (see, for example, Graves and Whinston 1970). We state it here for the sake of completeness.

**6.1. LEMMA.** *Let  $f: S_n \rightarrow \mathbb{R}$  be a function defined by (1.1.1) for some  $n \times n$  matrices  $A = (a_{ij})$ , and  $B = (b_{ij})$ . Let us define*

$$\begin{aligned} \alpha_1 &= \sum_{1 \leq i \neq j \leq n} a_{ij}, & \alpha_2 &= \sum_{i=1}^n a_{ii}, & \text{and} \\ \beta_1 &= \sum_{1 \leq i \neq j \leq n} b_{ij}, & \beta_2 &= \sum_{i=1}^n b_{ii}. \end{aligned}$$

Then,

$$\bar{f} = \frac{\alpha_1\beta_1}{n(n-1)} + \frac{\alpha_2\beta_2}{n}.$$

Similarly, if  $f$  is defined by (1.2.1) for some tensor  $C = \{c_{kl}^{ij}\}$ ,  $1 \leq i, j, k, l \leq n$ , then

$$\bar{f} = \frac{1}{n(n-1)} \sum_{1 \leq i \neq j \leq n} \sum_{1 \leq k \neq l \leq n} c_{kl}^{ij} + \frac{1}{n} \sum_{1 \leq i, l \leq n} c_{ll}^{ii}.$$

PROOF. We prove the first part only, because the proof of the the second part is completely similar. Let us choose a pair of indices  $1 \leq i \neq j \leq n$ . Then, as  $\sigma$  ranges over the symmetric group  $S_n$ , the ordered pair  $(\sigma(i), \sigma(j))$  ranges over all ordered pairs  $(k, l)$  with  $1 \leq k \neq l \leq n$  and each such a pair  $(k, l)$  appears  $(n-2)!$  times. Similarly, for each index  $1 \leq i \leq n$ , the index  $\sigma(i)$  ranges over the set  $\{1, \dots, n\}$  and each  $j \in \{1, \dots, n\}$  appears  $(n-1)!$  times. Therefore,

$$\begin{aligned} \bar{f} &= \frac{1}{n!} \sum_{\sigma \in S_n} \sum_{i, j=1}^n b_{\sigma(i)\sigma(j)} a_{ij} = \sum_{i, j=1}^n \left( a_{ij} \frac{1}{n!} \sum_{\sigma \in S_n} b_{\sigma(i)\sigma(j)} \right) \\ &= \frac{1}{n(n-1)} \sum_{i \neq j} a_{ij} \beta_1 + \frac{1}{n} \sum_{i=1}^n a_{ii} \beta_2 = \frac{\alpha_1 \beta_1}{n(n-1)} + \frac{\alpha_2 \beta_2}{n} \end{aligned}$$

and the proof follows.  $\square$

6.2. REMARK. Suppose that  $f(\sigma) = \langle B, \sigma(A) \rangle$  for some matrices  $A$  and  $B$  and all  $\sigma \in S_n$  is the objective function in the QAP (1.1), and suppose that the maximum value of  $f$  is attained at a permutation  $\tau$ . Let  $A_1 = \tau(A)$  and  $f_1(\sigma) = \langle B, \sigma(A_1) \rangle$ . Then  $f_1(\sigma) = f(\sigma\tau)$ , hence the maximum value of  $f_1$  is attained at the identity permutation  $\varepsilon$  and the distribution of values of  $f$  and  $f_1$  is the same. We observe that if  $A$  is symmetric, then  $A_1$  is also symmetric, and if  $A$  has constant row and column sums and a constant diagonal, then so does  $A_1$  (see also §7). Hence, as far as the distribution of values of  $f$  is concerned, without loss of generality we may assume that the maximum of  $f$  is attained at the identity permutation  $\varepsilon$ . The same is true for functions in the generalized problem (1.2).

Next, we introduce our main tool.

6.3. DEFINITION. Let  $f: S_n \rightarrow \mathbb{R}$  be a function. Let us define a function  $g: S_n \rightarrow \mathbb{R}$  by

$$g(\sigma) = \frac{1}{n!} \sum_{\omega \in S_n} f(\omega^{-1} \sigma \omega).$$

We call  $g$  the *central projection* of  $f$ .

It turns out that the central projection captures some important information regarding the distribution of values of a function.

6.4. LEMMA. Let  $f: S_n \rightarrow \mathbb{R}$  be a function and let  $g$  be the central projection of  $f$ . Then

(1) The averages of  $f$  and  $g$  over the  $k$ th sphere  $U(\varepsilon, k)$  around the identity permutation coincide:

$$\frac{1}{|U(\varepsilon, k)|} \sum_{\sigma \in U(\varepsilon, k)} f(\sigma) = \frac{1}{|U(\varepsilon, k)|} \sum_{\sigma \in U(\varepsilon, k)} g(\sigma).$$

(2) The average values of  $f$  and  $g$  on the symmetric group coincide:  $\bar{f} = \bar{g}$ .

(3) Suppose that  $f(\varepsilon) \geq f(\sigma)$  for all  $\sigma \in S_n$ . Then,  $g(\varepsilon) \geq g(\sigma)$  for all  $\sigma \in S_n$ .

PROOF. We observe that  $\sigma \in U(\varepsilon, k)$  if and only if  $\sigma$  has exactly  $k$  fixed points. Hence, for any fixed  $\omega \in S_n$ , the permutation  $\omega^{-1} \sigma \omega$  ranges over  $U(\varepsilon, k)$  as  $\sigma$  ranges over  $U(\varepsilon, k)$ . Hence,

$$\begin{aligned} \frac{1}{|U(\varepsilon, k)|} \sum_{\sigma \in U(\varepsilon, k)} g(\sigma) &= \frac{1}{|U(\varepsilon, k)|} \sum_{\sigma \in U(\varepsilon, k)} \left( \frac{1}{n!} \sum_{\omega \in S_n} f(\omega^{-1} \sigma \omega) \right) \\ &= \frac{1}{n!} \sum_{\omega \in S_n} \left( \frac{1}{|U(\varepsilon, k)|} \sum_{\sigma \in U(\varepsilon, k)} f(\omega^{-1} \sigma \omega) \right) = \frac{1}{|U(\varepsilon, k)|} \sum_{\sigma \in U(\varepsilon, k)} f(\sigma) \end{aligned}$$

and (1) is proven. Part (2) follows from (1). To prove (3), we note that  $\omega^{-1} \varepsilon \omega = \varepsilon$  for all  $\omega \in S_n$  and hence,  $g(\varepsilon) = f(\varepsilon)$ . Moreover, for any  $\sigma \in S_n$ ,

$$g(\sigma) = \frac{1}{n!} \sum_{\omega \in S_n} f(\omega^{-1} \sigma \omega) \leq \frac{1}{n!} \sum_{\omega \in S_n} f(\varepsilon) = g(\varepsilon). \quad \square$$

We will rely on a Markov type estimate, which asserts, roughly, that a function with a sufficiently large average takes sufficiently large values sufficiently often.

**6.5. LEMMA.** *Let  $X$  be a finite set and let  $f: X \rightarrow \mathbb{R}$  be a function. Suppose that  $f(x) \leq 1$  for all  $x \in X$  and that*

$$\frac{1}{|X|} \sum_{x \in X} f(x) \geq \beta \quad \text{for some } \beta > 0.$$

Then, for any  $0 < \gamma < 1$ , we have

$$|\{x \in X : f(x) \geq \beta\gamma\}| \geq \beta(1 - \gamma)|X|.$$

**PROOF.** We have

$$\begin{aligned} \beta &\leq \frac{1}{|X|} \sum_{x \in X} f(x) = \frac{1}{|X|} \sum_{x: f(x) < \beta\gamma} f(x) + \frac{1}{|X|} \sum_{x: f(x) \geq \beta\gamma} f(x) \\ &\leq \beta\gamma + \frac{|\{x : f(x) \geq \beta\gamma\}|}{|X|}. \end{aligned}$$

Hence,

$$|\{x : f(x) \geq \beta\gamma\}| \geq \beta(1 - \gamma)|X|. \quad \square$$

We need some facts about the structure of the symmetric group  $S_n$  (see, for example, Fulton and Harris 1991).

**6.6. The conjugacy classes of  $S_n$ .** Let us fix a permutation  $\rho \in S_n$ . As  $\omega$  ranges over the symmetric group  $S_n$ , the permutation  $\omega^{-1}\rho\omega$  ranges over the conjugacy class  $X(\rho)$  of  $\rho$ , that is the set of permutations that have the same cycle structure as  $\rho$ .

We will be using the following facts.

**6.6.1. Central projections and conjugacy classes.** If  $f: S_n \rightarrow \mathbb{R}$  is a function and  $g: S_n \rightarrow \mathbb{R}$  its central projection, then

$$g(\rho) = \frac{1}{|X(\rho)|} \sum_{\sigma \in X(\rho)} f(\sigma).$$

If  $X \subset S_n$  is a set which splits into a union of conjugacy classes  $X(\rho_i): i \in I$ , and for each such a class we have

$$\frac{1}{|X(\rho_i)|} \sum_{\sigma \in X(\rho_i)} f(\sigma) \geq \alpha$$

for some number  $\alpha$ , then

$$\frac{1}{|X|} \sum_{\sigma \in X} f(\sigma) \geq \alpha.$$

**6.6.2. Permutations with no fixed points and 2-cycles.** Let us fix some positive integers  $c_i: i = 1, \dots, m$  and let  $a_n$  be the number of permutations in  $S_n$  that have no cycles of length  $c_i$  for  $1 \leq i \leq m$ . The exponential generating function for  $a_n$  is given by

$$\sum_{n=0}^{\infty} \frac{a_n}{n!} x^n = \frac{1}{1-x} \exp \left\{ - \sum_{i=1}^m \frac{x^{c_i}}{c_i} \right\},$$

where we agree that  $a_0 = 1$ . (See, for example, pp. 170–173 of Goulden and Jackson 1983.) It follows that the number of permutations  $\sigma \in S_n$  without fixed points (known

as *derangements*) is asymptotically  $e^{-1}n!$ . More precisely, it is equal to  $d_n n!$  see (1.4). Similarly, the number of derangements without 2-cycles is asymptotically  $e^{-3/2}n!$ . We will use that the first number exceeds  $n!/3$  and the second number exceeds  $n!/5$  for  $n \geq 3$ .

We recall in (1.4) that  $p(\sigma)$  is the number of fixed points and  $t(\sigma)$  is the number of 2-cycles of a permutation  $\sigma \in S_n$ . We will need a couple of technical results.

6.7. LEMMA. *We have*

$$\frac{1}{|U(\varepsilon, k)|} \sum_{\sigma \in U(\varepsilon, k)} t(\sigma) = \frac{1}{2} \nu(n-k),$$

where  $\nu$  is the function of (1.4) and  $\varepsilon$  is the identity permutation.

PROOF. For a pair of indices  $1 \leq i < j \leq n$ , let

$$t_{ij}(\sigma) = \begin{cases} 1 & \text{if } \sigma(i) = j \text{ and } \sigma(j) = i, \\ 0 & \text{otherwise.} \end{cases}$$

Then,  $t(\sigma) = \sum_{i < j} t_{ij}(\sigma)$ . Let us compute the average of  $t_{ij}(\sigma)$  over  $U(\varepsilon, k)$ . To choose a permutation  $\sigma \in U(\varepsilon, k)$ , one has to choose  $k$  fixed points in  $\binom{n}{k}$  ways and then a derangement on the remaining  $n-k$  symbols in  $d_{n-k}(n-k)!$  ways. Hence,  $|U(\varepsilon, k)| = d_{n-k} n! / k!$ . To choose a permutation  $\sigma \in U(\varepsilon, k)$  where  $(ij)$  is a 2-cycle, one has to choose  $k$  fixed points in  $\binom{n-2}{k}$  ways and then a derangement on  $n-k-2$  symbols. Hence, the total number of permutations  $\sigma \in U(\varepsilon, k)$  with  $t_{ij}(\sigma) = 1$  is

$$\binom{n-2}{k} d_{n-k-2} (n-k-2)! = \frac{(n-2)! d_{n-k-2}}{k!}.$$

Thus, for all pairs  $i < j$  we have

$$\frac{1}{|U(\varepsilon, k)|} \sum_{\sigma \in U(\varepsilon, k)} t_{ij}(\sigma) = \frac{(n-2)! d_{n-k-2}}{n! d_{n-k}} = \frac{\nu(n-k)}{n(n-1)}$$

and

$$\frac{1}{|U(\varepsilon, k)|} \sum_{\sigma \in U(\varepsilon, k)} t(\sigma) = \sum_{i < j} \frac{1}{|U(\varepsilon, k)|} \sum_{\sigma \in U(\varepsilon, k)} t_{ij}(\sigma) = \binom{n}{2} \frac{\nu(n-k)}{n(n-1)} = \frac{\nu(n-k)}{2}. \quad \square$$

6.8. LEMMA. *For a permutation  $\sigma \in S_n$ ,  $\sigma \neq \varepsilon$ , let  $a_\sigma \in \mathbb{R}^2$  be the point*

$$a_\sigma = \left( p(\sigma), \frac{2t(\sigma)}{n-p(\sigma)} \right).$$

*Let  $P = \text{conv}\{a_\sigma : \sigma \neq \varepsilon\}$  be the convex hull of all such points  $a_\sigma$ .*

*If  $n$  is even, the extreme points of  $P$  are*

$$(0, 0), \quad (n-3, 0), \quad (n-2, 1), \quad \text{and} \quad (0, 1).$$

*If  $n$  is odd, the extreme points of  $P$  are*

$$(0, 0), \quad (n-3, 0), \quad (n-2, 1), \quad (0, (n-3)/n), \quad \text{and} \quad (1, 1).$$

PROOF. The set of all possible values  $(p(\sigma), t(\sigma))$ , where  $\sigma \neq \varepsilon$ , consists of all pairs of nonnegative integers  $(p, t)$  such that  $p \leq n-2$ ,  $2t \leq n$  and, additionally,  $p+2t \leq n-3$  or  $p+2t = n$ . To find the extreme points of the set of feasible points  $(p, 2t/(n-p))$ , we choose a generic vector  $(\gamma_1, \gamma_2)$  and investigate for which values of  $p$  and  $t$  the maximum of

$$\gamma_1 p + \gamma_2 \frac{2t}{n-p}$$

is attained.

Clearly, we can assume that  $\gamma_2 \neq 0$ . If  $\gamma_2 < 0$ , then we should choose the smallest possible  $t$  which would be  $t = 0$  unless  $p = n-2$  when we have to choose  $t = 1$ . Depending on the sign of  $\gamma_1$ , this produces the following pairs:

$$(p, t) = \{(0, 0), (n-3, 0), (n-2, 1)\}.$$

If  $\gamma_2 > 0$ , then the largest possible value of  $2t/(n-p)$  is 1. If  $\gamma_1 > 0$ , this produces the (already included) point

$$(p, t) = (n-2, 1).$$

If  $\gamma_1 < 0$ , we get

$$(p, t) = (0, n/2) \quad \text{for even } n$$

and

$$(p, t) = \{(0, (n-3)/2), (1, (n-1)/2)\} \quad \text{for odd } n.$$

Summarizing, the extreme points of  $P$  are

$$(0, 0), (n-3, 0), (n-2, 1), (0, 1) \quad \text{for even } n$$

and

$$(0, 0), (n-3, 0), (n-2, 1), (0, (n-3)/n), (1, 1) \quad \text{for odd } n$$

as claimed.  $\square$

**7. Action of the symmetric group in the space of matrices.** The crucial observation for our approach is that the vector space of all central projections  $g$  of functions  $f$  defined by (1.1.1) or (1.2.1) is 4-, 3-, or 2-dimensional depending on whether we consider the general case, the cases of §§3 and 4, or the special case of §2. If we require, additionally, that  $\bar{f} = 0$ , then the dimensions drop by 1 to 3, 2 and 1, respectively. This fact is explained by the representation theory of the symmetric group (see, for example, Fulton and Harris 1991). In this section, we review some facts that we need. Our notation is inspired by the generally accepted notation of representation theory.

We describe some important invariant subspaces of the action of  $S_n$  in the space of  $n \times n$  matrices  $\text{Mat}_n$  by simultaneous permutations of rows and columns. We recall that  $n \geq 4$ .

**7.1. Subspace  $L_n$ .** Let  $L_n^1$  be the space of constant matrices  $A$ :

$$a_{ij} = \alpha \quad \text{for some } \alpha \text{ and all } 1 \leq i, j \leq n.$$

Let  $L_n^2$  be the subspace of scalar matrices  $A$ :

$$a_{ij} = \begin{cases} \alpha & \text{if } i = j \\ 0 & \text{if } i \neq j \end{cases} \quad \text{for some } \alpha.$$

Finally, Let  $L_n = L_n^1 + L_n^2$ . One can observe that  $\dim L_n = 2$  and that  $L_n$  is the subspace of all matrices that remain fixed under the action of  $S_n$ .

**7.2. Subspace  $L_{n-1,1}$ .** Let  $L_{n-1,1}^1$  be the subspace of matrices with identical rows and such that the sum of entries in each row is 0:

$$A = \begin{pmatrix} \alpha_1 & \alpha_2 & \dots & \alpha_n \\ \alpha_1 & \alpha_2 & \dots & \alpha_n \\ \dots & \dots & \dots & \dots \\ \alpha_1 & \alpha_2 & \dots & \alpha_n \end{pmatrix} \quad \text{where } \alpha_1 + \dots + \alpha_n = 0.$$

Similarly, let  $L_{n-1,1}^2$  be the subspace of matrices with identical columns and such that the sum of entries in each column is 0:

$$A = \begin{pmatrix} \alpha_1 & \alpha_1 & \dots & \alpha_1 \\ \alpha_2 & \alpha_2 & \dots & \alpha_2 \\ \dots & \dots & \dots & \dots \\ \alpha_n & \alpha_n & \dots & \alpha_n \end{pmatrix} \quad \text{where } \alpha_1 + \dots + \alpha_n = 0.$$

Finally, let  $L_{n-1,1}^3$  be the subspace of diagonal matrices whose diagonal entries sum to zero:

$$A = \begin{pmatrix} \alpha_1 & 0 & \dots & 0 & 0 \\ 0 & \alpha_2 & \dots & 0 & 0 \\ \dots & \dots & \dots & \dots & \dots \\ 0 & 0 & \dots & 0 & \alpha_n \end{pmatrix} \quad \text{where } \alpha_1 + \dots + \alpha_n = 0.$$

Let  $L_{n-1,1} = L_{n-1,1}^1 + L_{n-1,1}^2 + L_{n-1,1}^3$ . One can check that the dimension of each of  $L_{n-1,1}^1$ ,  $L_{n-1,1}^2$ , and  $L_{n-1,1}^3$  is  $n-1$  and that  $\dim L_{n-1,1} = 3n-3$ . Moreover, the subspaces  $L_{n-1,1}^1$ ,  $L_{n-1,1}^2$ , and  $L_{n-1,1}^3$  do not contain nontrivial invariant subspaces. The action of  $S_n$  in  $L_{n-1,1}$ , although nontrivial, is not very complicated. One can show that if  $A \in L_{n-1,1} + L_n$ , then problem 1.1 of optimizing  $f(\sigma)$  reduces to the LAP.

**7.3. Subspace  $L_{n-2,2}$ .** Let us define  $L_{n-2,2}$  as the subspace of all *symmetric* matrices  $A$  with row and column sums equal to 0 and zero diagonal:

$$\begin{aligned} a_{ij} &= a_{ji} && \text{for all } 1 \leq i, j \leq n, \\ \sum_{i=1}^n a_{ij} &= 0 && \text{for all } j = 1, \dots, n, \\ \sum_{j=1}^n a_{ij} &= 0 && \text{for all } i = 1, \dots, n, \quad \text{and} \\ a_{ii} &= 0 && \text{for all } i = 1, \dots, n. \end{aligned}$$

One can check that  $L_{n-2,2}$  is an invariant subspace and that  $\dim L_{n-2,2} = (n^2 - 3n)/2$ . Besides,  $L_{n-2,2}$  contains no nontrivial invariant subspaces.

**7.4. Subspace  $L_{n-2,1,1}$ .** Let us define  $L_{n-2,1,1}$  as the subset of all *skew symmetric* matrices  $A$  with row and column sums equal to 0:

$$\begin{aligned} a_{ij} &= -a_{ji} && \text{for all } 1 \leq i, j \leq n, \\ \sum_{i=1}^n a_{ij} &= 0 && \text{for all } j = 1, \dots, n, \quad \text{and} \\ \sum_{j=1}^n a_{ij} &= 0 && \text{for all } i = 1, \dots, n. \end{aligned}$$



One can check that  $L_{n-2,1,1}$  is an invariant subspace and that  $\dim L_{n-2,1,1} = (n^2 - 3n)/2 + 1$ . Similarly,  $L_{n-2,1,1}$  contains no nontrivial invariant subspaces.

One can check that  $\text{Mat}_n = L_n + L_{n-1,1} + L_{n-2,2} + L_{n-2,1,1}$ . The importance of the subspaces (7.1)–(7.4) is explained by the fact that they are the *isotypical components* of the irreducible representations of the symmetric group in the space of matrices.

To state the main result of this section, we recall the definitions of the central projection (see Definition 6.3) and of the functions  $p$  and  $t$  (see 1.4).

**7.5. PROPOSITION.** *For  $n \times n$  matrices  $A$  and  $B$ , where  $n \geq 4$ , let  $f: S_n \rightarrow \mathbb{R}$  be the function defined by (1.1.1) and let  $g: S_n \rightarrow \mathbb{R}$  be the central projection of  $f$ .*

(1) *If  $A \in L_n$ , then  $g$  is a scalar multiple of the constant function*

$$\chi_n(\sigma) = 1 \quad \text{for all } \sigma \in S_n.$$

(2) *If  $A \in L_{n-1,1}$ , then  $g$  is a scalar multiple of the function*

$$\chi_{n-1,1}(\sigma) = p(\sigma) - 1 \quad \text{for all } \sigma \in S_n.$$

(3) *If  $A \in L_{n-2,2}$ , then  $g$  is a scalar multiple of the function*

$$\chi_{n-2,2}(\sigma) = t(\sigma) + \frac{1}{2}p^2(\sigma) - \frac{3}{2}p(\sigma) \quad \text{for all } \sigma \in S_n.$$

(4) *If  $A \in L_{n-2,1,1}$ , then  $g$  is a scalar multiple of the function*

$$\chi_{n-2,1,1}(\sigma) = \frac{1}{2}p^2(\sigma) - \frac{3}{2}p(\sigma) - t(\sigma) + 1 \quad \text{for all } \sigma \in S_n.$$

Proposition 7.5 follows from the representation theory of the symmetric group (see, for example, Part 1 of Fulton and Harris 1991). The set of all functions  $f: S_n \rightarrow \mathbb{R}$  is identified with the (real) group algebra of the symmetric group. The center of the group algebra is spanned by the characters of the irreducible representations of  $S_n$ . The basic fact that we are using here is the following: if  $f$  is a matrix element in an irreducible representation of the group, then the central projection must be a scalar multiple of the character of that representation. The functions  $\chi_n, \chi_{n-1,1}, \chi_{n-2,2}$ , and  $\chi_{n-2,1,1}$  are the characters of corresponding irreducible representations of  $S_n$  for  $n \geq 4$  (see Lecture 4 of Fulton and Harris 1991). The irreducible characters are linearly independent, and, moreover, orthonormal. In particular, we will use that

$$\sum_{\sigma \in S_n} \chi_{n-1,1}(\sigma) = \sum_{\sigma \in S_n} \chi_{n-2,2}(\sigma) = \sum_{\sigma \in S_n} \chi_{n-2,1,1}(\sigma) = 0.$$

Hence, the average value of all but the trivial character  $\chi_n$  is 0.

**7.6. REMARK.** We note that the functions  $p, p^2$ , and  $t$  are objective functions of type (1.2.1) in some generalized QAP (see 1.2). Indeed, to obtain  $p$  we choose  $c_{ii}^{ii} = 1$  for all  $i = 1, \dots, n$  to be the only nonzero entries of  $C$ . To obtain  $p^2$ , we choose  $c_{ij}^{ij} = 1$  for  $1 \leq i, j \leq n$  to be the only nonzero entries of  $C$ . To obtain  $t$ , we choose  $c_{ji}^{ij} = 1$  for all  $1 \leq i < j \leq n$  to be the only nonzero entries of  $C$ . Consequently, the characters  $\chi_n, \chi_{n-1,1}, \chi_{n-2,2}$ , and  $\chi_{n-2,1,1}$  are objective functions of type (1.2.1).

**8. The bullseye case: Proofs.** In this section, we prove Theorem 2.1 and Theorem 2.3. The proof is based on the observation that  $A$  satisfies the conditions of §2 if and only if  $A \in L_n + L_{n-2,2}$  (see §7).

PROOF OF THEOREM 2.1. Without loss of generality, we may assume that the maximum of  $f_0(\sigma)$  is attained at the identity permutation  $\varepsilon$  (see Remark 6.2). Excluding the noninteresting case of  $f_0 \equiv 0$ , by scaling  $f$ , if necessary, we can assume that  $f_0(\varepsilon) = 1$ . Let  $g$  be the central projection of  $f_0$ . Then, by Parts (2) and (3) of Lemma 6.4, we have  $\bar{g} = 0$  and  $1 = g(\varepsilon) \geq g(\sigma)$  for all  $\sigma \in S_n$ . Moreover, because  $A \in L_n + L_{n-2,2}$ , by Parts (1) and (3) of Proposition 7.5,  $g$  must be a linear combination of the constant function  $\chi_n$  and  $\chi_{n-2,2}$ . Because  $\bar{g} = 0$ ,  $g$  should be proportional to  $\chi_{n-2,2}$ , and because  $g(\varepsilon) = 1$ , we have

$$g = \frac{2}{n^2 - 3n} \chi_{n-2,2} = \frac{2t + p^2 - 3p}{n^2 - 3n}.$$

Now,  $\sigma \in U(\varepsilon, k)$  if and only if  $p(\sigma) = k$ . Applying Part (1) of Lemma 6.4 and Lemma 6.7, we get

$$\frac{1}{|U(\varepsilon, k)|} \sum_{\sigma \in U(\varepsilon, k)} f_0(\sigma) = \frac{1}{|U(\varepsilon, k)|} \sum_{\sigma \in U(\varepsilon, k)} g(\sigma) = \frac{k^2 - 3k + \nu(n-k)}{n^2 - 3n}$$

and the proof follows.  $\square$

PROOF OF THEOREM 2.3. As in the proof of Theorem 2.1, we assume that the maximum value of  $f_0$  is equal to 1.

Let us estimate the cardinality  $|U(\tau, k)| = |U(\varepsilon, k)|$ . Because  $\sigma \in U(\varepsilon, k)$  if and only if  $\sigma$  has  $k$  fixed points, to choose a  $\sigma \in U(\varepsilon, k)$  one has to choose  $k$  points in  $\binom{n}{k}$  ways and then choose a derangement on the remaining  $n-k$  points. Using Subsection 6.6.2, we get

$$|U(\tau, k)| \geq \binom{n}{k} (n-k)!/3 = \frac{n!}{3k!}.$$

Applying Lemma 6.5 with  $\beta = \beta(n, k)$  and  $X = U(\tau, k)$ , from Theorem 2.1 we conclude that

$$\begin{aligned} \mathbf{P}\{\sigma \in S_n: f_0(\sigma) \geq \gamma\beta(n, k)\} &\geq \frac{(1-\gamma)\beta(n, k)|U(\tau, k)|}{n!} \\ &\geq \frac{(1-\gamma)\beta(n, k)}{3k!}. \quad \square \end{aligned}$$

**9. The pure case: Proofs.** In this section, we prove Theorem 3.1 and Theorem 3.2. We observe that  $A$  satisfies the conditions of §3 if and only if  $A \in L_n + L_{n-2,1,1} + L_{n-2,2}$  (see §7). As in §8, the  $L_n$  component contributes just a constant to  $f$ . Because the  $L_{n-1,1}$  component attributed to the LAP (see §7.2) is absent, we call this case “pure.”

We choose a more convenient basis  $g_1$  and  $g_2$  in the vector space spanned by  $\chi_{n-2,2}$  and  $\chi_{n-2,1,1}$ , namely

$$g_1 = \chi_{n-2,2} + \chi_{n-2,1,1} = p^2 - 3p + 1 \quad \text{and} \quad g_2 = \chi_{n-2,1,1} - \chi_{n-2,2} = 1 - 2t.$$

9.1. DEFINITION. Let  $K_p$  (where  $p$  stands for “pure”) be the set of all functions  $g: S_n \rightarrow \mathbb{R}$  such that  $g \in \text{span}\{g_1, g_2\}$ , where  $g_1 = p^2 - 3p + 1$ ,  $g_2 = 1 - 2t$ , and  $g(\varepsilon) \geq g(\sigma)$  for all  $\sigma \in S_n$ , where  $\varepsilon$  is the identity permutation. We call  $K_p$  the *central (pure) cone*.

Identifying  $\text{span}\{g_1, g_2\}$  with two-dimensional vector space  $\mathbb{R}^2$  (plane), we see that the conditions  $g(\varepsilon) \geq g(\sigma)$  define the central cone  $K_p$  as a convex cone in  $\mathbb{R}^2$ . Our goal is to find the extreme rays  $r_1$  and  $r_2$  of  $K_p$ , so that every function  $g \in K_p$  can be written as a nonnegative linear combination of  $r_1$  and  $r_2$ .

9.2. LEMMA. For  $n \geq 4$ , let us define the functions  $r_1, r_{2e}$ , and  $r_{2o}: S_n \rightarrow \mathbb{R}$  by

$$\begin{aligned} r_1 &= 1 - 2t, \\ r_{2e} &= \frac{p^2 - 3p - n - 6t + 2tn + 4}{n^2 - 4n + 4}, \quad \text{and} \\ r_{2o} &= \frac{p^2 - 3p - n - 4t + 2tn + 3}{n^2 - 4n + 3}. \end{aligned}$$

Then,

- (1) If  $n$  is even, then  $K_p$  is a 2-dimensional convex cone with the extreme rays spanned by  $r_1$  and  $r_{2e}$ .
- (2) If  $n$  is odd, then  $K_p$  is a 2-dimensional convex cone with the extreme rays spanned by  $r_1$  and  $r_{2o}$ . Cone  $K_p$  contains  $r_{2e}$ .
- (3) If  $\varepsilon \in S_n$  is the identity, then

$$r_1(\varepsilon) = r_{2e}(\varepsilon) = r_{2o}(\varepsilon) = 1.$$

PROOF. A function  $g \in K_p$  can be written as a linear combination  $g = \alpha_1 g_1 + \alpha_2 g_2$ . Because  $p(\varepsilon) = n$  and  $t(\varepsilon) = 0$ , we have  $g(\varepsilon) = \alpha_1(n^2 - 3n + 1) + \alpha_2$ . Therefore, the inequalities  $g(\varepsilon) \geq g(\sigma)$  can be written as

$$\alpha_1(n^2 - 3n + 1) + \alpha_2 \geq \alpha_1(p^2(\sigma) - 3p(\sigma) + 1) + \alpha_2(1 - 2t(\sigma)),$$

which, for  $\sigma \neq \varepsilon$  is equivalent to

$$\alpha_1(n + p(\sigma) - 3) + \alpha_2 \frac{2t(\sigma)}{n - p(\sigma)} \geq 0.$$

Using Lemma 6.8, we conclude that for even  $n$ , the system is equivalent to

$$(9.2.1) \quad \begin{aligned} \alpha_1 &\geq 0, \\ (n-3)\alpha_1 + \alpha_2 &\geq 0, \end{aligned}$$

and for odd  $n$ , the system is equivalent to

$$(9.2.2) \quad \begin{aligned} \alpha_1 &\geq 0, \\ (n-2)\alpha_1 + \alpha_2 &\geq 0. \end{aligned}$$

Consequently, every solution  $(\alpha_1, \alpha_2)$  of (9.2.1) is a nonnegative linear combination of  $(0, 1)$  and  $(1, 3 - n)$ , and every solution of (9.2.2) is a nonnegative linear combination of  $(0, 1)$  and  $(1, 2 - n)$ .

The functions  $r_1, r_{2e}$ , and  $r_{2o}$  are obtained from  $g_2, g_1 + (3 - n)g_2$  and  $g_1 + (2 - n)g_2$  respectively by scaling so that the value at the identity becomes equal to 1.

Because every solution of (9.2.1) is a solution of (9.2.2), we conclude that  $r_{2e} \in K_p$  for odd  $n$  as well.  $\square$

9.3. REMARK. We observe that  $r_1$  has the damped oscillator distribution corresponding to the case of  $\gamma_1 = 1$  in Theorem 3.1, whereas  $r_{2e}$  and  $r_{2o}$  both have the bullseye distribution corresponding to the case of  $\gamma_2 = 1$  in Theorem 3.1. If  $n$  is even, then  $r_{2o} \notin K_p$ ; see Figure 5. Indeed, if  $\sigma$  is a product of  $n/2$  commuting 2-cycles, so that  $p(\sigma) = 0$  and  $t(\sigma) = n/2$ , then  $r_{2o}(\sigma) = (n^2 - 3n + 3)/(n^2 - 4n + 3) > 1 = r_{2o}(\varepsilon)$ .

PROOF OF THEOREM 3.1. We proceed as in the proof of Theorem 2.1 (§8) with some modifications. Without loss of generality, we assume that the maximum of  $f_0(\sigma)$  is attained on the identity permutation  $\varepsilon$  and that  $f_0(\varepsilon) = 1$ . Let  $g$  be the central projection of  $f_0$ . Because  $A \in L_n + L_{n-2,2} + L_{n-2,1,1}$ , by Parts (1), (3), and (4) of Proposition 7.5,  $g$  is a linear

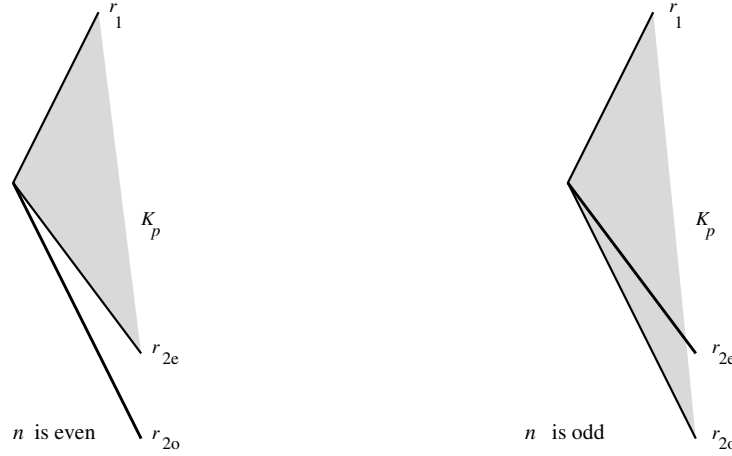


FIGURE 5. The central (pure) cone.

combination of  $\chi_n$ ,  $\chi_{n-2,2}$ , and  $\chi_{n-2,1,1}$ . By Part (2) of Lemma 6.4, we have  $\bar{g} = \bar{f}_0 = 0$ , so  $g$  is a linear combination of  $\chi_{n-2,2}$  and  $\chi_{n-2,1,1}$  alone. Moreover, by Part (3) of Lemma 6.4, we have  $1 = g(\varepsilon) \geq g(\sigma)$  for all  $\sigma \in S_n$ . Hence,  $g$  lies in the central cone  $K_p$ . Applying Lemma 9.2, we conclude that  $g$  must be a convex combination of  $r_1$  and  $r_{2e}$  for  $n$  even and a convex combination of  $r_1$  and  $r_{2o}$  for  $n$  odd. Applying Part (1) of Lemma 6.4, we can replace the average of  $f_0$  over the set  $U(\varepsilon, k)$  by the average of  $g$  over  $U(\varepsilon, k)$ . The proof now follows by Lemma 6.7 and the observation that  $\sigma \in U(\varepsilon, k)$  if and only if  $p(\sigma) = k$ .  $\square$

To prove Theorem 3.2, we need need one preliminary result.

9.4. LEMMA. *Let  $g$  be a linear combination of  $g_1 = \chi_{n-2,2} + \chi_{n-2,1,1} = p^2 - 3p + 1$  and  $g_2 = \chi_{n-2,1,1} - \chi_{n-2,2} = 1 - 2t$  such that  $g(\varepsilon) = 1$ . For  $3 \leq k \leq n-1$ , let  $\sigma_k$  be a permutation such that  $p(\sigma_k) = k$  and  $t(\sigma_k) = 0$ , and let  $\theta_k$  be a permutation such that  $p(\theta_k) = k$  and  $t(\theta_k) = 1$ . Then,*

$$\max\{g(\sigma_k), g(\theta_k)\} \geq \frac{k^2 - 3k + 1}{n^2 - 3n + 1}.$$

PROOF. Because  $g(\varepsilon) = 1$ ,  $g_1(\varepsilon) = n^2 - 3n + 1$ , and  $g_2(\varepsilon) = 1$ , we can write

$$g = \alpha_1 \frac{p^2 - 3p + 1}{n^2 - 3n + 1} + \alpha_2 (1 - 2t)$$

for some  $\alpha_1$  and  $\alpha_2$  such that  $\alpha_1 + \alpha_2 = 1$ . Then,

$$\begin{aligned} g(\sigma_k) &= \alpha_1 \frac{k^2 - 3k + 1}{n^2 - 3n + 1} + \alpha_2 \quad \text{and} \\ g(\theta_k) &= \alpha_1 \frac{k^2 - 3k + 1}{n^2 - 3n + 1} - \alpha_2. \end{aligned}$$

We observe that  $g(\sigma_k)$  and  $g(\theta_k)$  are linear functions of  $\alpha_1$  and  $\alpha_2$ , and that for  $\alpha_1 = 1$  and  $\alpha_2 = 0$  we have

$$(9.4.1) \quad g(\sigma_k) = g(\theta_k) = \frac{k^2 - 3k + 1}{n^2 - 3n + 1}.$$

Let

$$\lambda_1 = \frac{n^2 - 3n + 1}{2(k^2 - 3k + 1)} + \frac{1}{2} \quad \text{and} \quad \lambda_2 = \frac{n^2 - 3n + 1}{2(k^2 - 3k + 1)} - \frac{1}{2}.$$

Then,  $\lambda_1, \lambda_2 > 0$  and

$$\lambda_1 g(\sigma_k) + \lambda_2 g(\theta_k) = \alpha_1 + \alpha_2 = 1.$$

Comparing this with (9.4.1) we conclude that there are no values  $\alpha_1$  and  $\alpha_2$  such that  $\alpha_1 + \alpha_2 = 1$  and

$$g(\sigma_k), g(\theta_k) < \frac{k^2 - 3k + 1}{n^2 - 3n + 1}. \quad \square$$

**PROOF OF THEOREM 3.2.** Without loss of generality, we may assume that the maximum value of  $f_0$  is attained at the identity permutation  $\varepsilon$  (see Remark 6.2). Excluding an obvious case of  $f_0 \equiv 0$ , by scaling  $f$ , if necessary, we may assume that  $f_0(\varepsilon) = 1$ . Let  $g$  be the central projection of  $f_0$ . As in the proof of Theorem 3.1, we deduce that  $g$  is a linear combination of  $\chi_{n-2,2}$  and  $\chi_{n-2,1,1}$ , and hence a linear combination of  $g_1$  and  $g_2$ , and that  $g(\varepsilon) = 1$ .

Let us choose  $3 \leq k \leq n-3$  and let  $X_k$  be the set of permutations  $\sigma$  such that  $p(\sigma) = k$  and  $t(\sigma) = 0$ , and let  $Y_k$  be the set of permutations  $\theta$  such that  $p(\theta) = k$  and  $t(\theta) = 1$ . To choose a permutation  $\sigma \in X_k$ , one has to choose  $k$  fixed points in  $\binom{n}{k}$  ways and then a derangement without 2-cycles on the remaining  $(n-k)$  points. Then, by (6.6.2),

$$|X_k| \geq \frac{1}{5} \binom{n}{k} (n-k)! = \frac{1}{5} \frac{n!}{k!}.$$

Similarly, to choose a permutation  $\theta \in Y_k$ , one has to choose a 2-cycle in  $\binom{n}{2}$  ways  $k$  fixed points in  $\binom{n-2}{k}$  ways, and a derangement without 2-cycles on the remaining  $(n-k-2)$  points. Then, by (6.6.2),

$$|Y_k| \geq \frac{1}{5} \binom{n}{2} \binom{n-2}{k} (n-k-2)! = \frac{n!}{10k!}.$$

Let us choose a permutation  $\sigma \in X_k$  and a permutation  $\theta \in Y_k$ , and let  $Z = X_k$  if  $g(\sigma) \geq g(\theta)$  and  $Z = Y_k$  otherwise. Then,

$$|Z| \geq \frac{n!}{10k!}$$

and by Lemma 9.4,

$$g(\sigma) \geq \frac{k^2 - 3k + 1}{n^2 - 3n + 1} \quad \text{for all } \sigma \in Z.$$

The set  $Z$  is a disjoint union of some conjugacy classes  $X(\rho)$  and for each  $X(\rho)$  by (6.6.1), we have

$$g(\rho) = \frac{1}{|X(\rho)|} \sum_{\sigma \in X(\rho)} f_0(\sigma) \geq \frac{k^2 - 3k + 1}{n^2 - 3n + 1}$$

and hence

$$\frac{1}{|Z|} \sum_{\sigma \in Z} f_0(\sigma) \geq \frac{k^2 - 3k + 1}{n^2 - 3n + 1}.$$

Applying Lemma 6.5 with  $X = Z$  and  $\beta = \beta(n, k)$ , we get that

$$\mathbf{P}\{\sigma \in S_n : f_0(\sigma) \geq \gamma \beta(n, k)\} \geq \frac{(1-\gamma)\beta(n, k)}{10k!}. \quad \square$$

To obtain Corollary 3.3 from Theorem 3.2, we fix any  $\gamma \in (0, 1)$  (say  $\gamma = 1/2$ ) and choose  $k = O(\sqrt{\alpha})$  in Part (1), and  $k = O(n^{1-\epsilon/2})$  in Part (2).

**9.5. REMARK.** It follows from the proof that we are able to choose the required number of “good” permutations among the permutations whose distance to the optimal permutation  $\tau$  is  $n-k$ .

**10. The symmetric case: Proofs.** In this case,  $A \in L_n + L_{n-1,1} + L_{n-2,2}$  (see §7). As in §§8 and 9, the  $L_n$  component contributes just a constant to  $f$ . We choose a more convenient basis  $g_1$  and  $g_2$  in the vector space spanned by  $\chi_{n-1,1}$  and  $\chi_{n-2,2}$ , namely

$$g_1 = \chi_{n-1,1} = p - 1 \quad \text{and} \quad g_2 = 2\chi_{n-2,2} + 3\chi_{n-1,1} = p^2 + 2t - 3,$$

where  $p(\sigma)$  is the number of fixed points of  $\sigma$  and  $t(\sigma)$  is the number of 2-cycles in  $\sigma$ .

10.1. DEFINITION. Let  $K_s$  (where  $s$  stands for ‘‘symmetric’’) be the set of all functions  $g: S_n \rightarrow \mathbb{R}$  such that  $g \in \text{span}\{g_1, g_2\}$ , where  $g_1 = p - 1$ ,  $g_2 = p^2 + 2t - 3$ , and  $g(\varepsilon) \geq g(\sigma)$  for all  $\sigma \in S_n$ , where  $\varepsilon$  is the identity permutation. We call  $K_s$  the *central (symmetric) cone*.

Identifying  $\text{span}\{g_1, g_2\}$  with 2-dimensional vector space  $\mathbb{R}^2$  (plane), we see that the conditions  $g(\varepsilon) \geq g(\sigma)$  define the central cone  $K_s$  as a convex cone in  $\mathbb{R}^2$ . Our goal is to find the extreme rays  $r_1$  and  $r_2$  of  $K_s$ , so that every function  $g \in K_s$  can be written as a nonnegative linear combination of  $r_1$  and  $r_2$ .

10.2. LEMMA. For  $n \geq 4$ , let us define the functions  $r_1, r_{2e}$ , and  $r_{2o}: S_n \rightarrow \mathbb{R}$  by

$$\begin{aligned} r_1 &= \frac{2np - 2n - p^2 - 3p - 2t + 6}{n^2 - 5n + 6}, \\ r_{2e} &= \frac{-np + n + p^2 + p + 2t - 4}{2n - 4}, \quad \text{and} \\ r_{2o} &= \frac{-n^2p + np^2 + n^2 + np + 2nt - 4n - 3p + 3}{2n^2 - 7n + 3}. \end{aligned}$$

Then,

- (1) If  $n$  is even, then  $K_s$  is a 2-dimensional convex cone with the extreme rays spanned by  $r_1$  and  $r_{2e}$ .
- (2) If  $n$  is odd, then  $K_s$  is a 2-dimensional convex cone with the extreme rays spanned by  $r_1$  and  $r_{2o}$ . Cone  $K_s$  contains  $r_{2e}$ .
- (3) If  $\varepsilon \in S_n$  is the identity, then

$$r_1(\varepsilon) = r_{2e}(\varepsilon) = r_{2o}(\varepsilon) = 1.$$

PROOF. A function  $g \in K_s$  can be written as a linear combination  $g = \alpha_1 g_1 + \alpha_2 g_2$ . Because  $p(\varepsilon) = n$  and  $t(\varepsilon) = 0$ , we have  $g(\varepsilon) = \alpha_1(n - 1) + \alpha_2(n^2 - 3)$ . Therefore, the inequalities  $g(\varepsilon) \geq g(\sigma)$  can be written as

$$\alpha_1(n - 1) + \alpha_2(n^2 - 3) \geq \alpha_1(p(\sigma) - 1) + \alpha_2(p^2(\sigma) + 2t(\sigma) - 3),$$

which, for  $\sigma \neq \varepsilon$ , is equivalent to

$$(10.2.1) \quad \alpha_1 + \alpha_2 \left( n + p(\sigma) - \frac{2t(\sigma)}{n - p(\sigma)} \right) \geq 0.$$

Applying Lemma 6.8, we observe that (10.2.1) is equivalent to the system of two inequalities:

$$\alpha_1 + (2n - 3)\alpha_2 \geq 0$$

and

$$\begin{aligned} \alpha_1 + (n - 1)\alpha_2 &\geq 0 && \text{if } n \text{ is even,} \\ n\alpha_1 + (n^2 - n + 3)\alpha_2 &\geq 0 && \text{if } n \text{ is odd.} \end{aligned}$$

Thus, every pair  $(\alpha_1, \alpha_2)$  satisfying (10.2.1) can be written as a nonnegative linear combination of  $(2n - 3, -1)$  and  $(1 - n, 1)$  when  $n$  is even, and  $(2n - 3, -1)$  and  $(-n^2 + n - 3, n)$  when  $n$  is odd.

The generators  $r_1, r_{2e}$ , and  $r_{2o}$  are obtained from  $(2n-3)g_1 - g_2$ ,  $(1-n)g_1 + g_2$ , and  $(-n^2 + n - 3)g_1 + ng_2$ , respectively by scaling so that the value at the identity becomes equal to 1.

It remains to check that  $r_{2e} \in K_s$  for  $n$  odd as well. Indeed, using that  $2t + p \leq n$  we have

$$\begin{aligned} (2n-4)(r_{2e} - 1) &= -n(p-1) + p(p+1) + 2t - 4 - 2n + 4 \\ &= -n(p+1) + p(p+1) + 2t \\ &\leq (p+1)(-n+p) + n - p = p(-n+p) \leq 0. \quad \square \end{aligned}$$

10.3. REMARK. We observe that  $r_1$  has the bullseye distribution corresponding to the case of  $\gamma_1 = 1$  in Theorem 4.1, whereas  $r_{2e}$  and  $r_{2o}$  both have the spike distribution corresponding to the case of  $\gamma_2 = 1$  in Theorem 4.1. One can observe that if  $n$  is even, then  $r_{2o} \notin K_s$ . Indeed, if  $\sigma$  is a product of  $n/2$  commuting 2-cycles, then  $r_{2o}(\sigma) = (2n^2 - 4n + 3)/(2n^2 - 7n + 3) > 1$ . The picture of  $K_s$  is very similar to that of  $K_p$  (see §9, Figure 5).

PROOF OF THEOREM 4.1. The proof is completely similar to the proof of Theorem 3.1 (see §9). We use Lemma 10.2 instead of Lemma 9.2.  $\square$

To prove Theorem 4.3, we need one preliminary result.

10.4. LEMMA. *Let  $g$  be a linear combination of  $g_1 = \chi_{n-1,1} = p-1$  and  $g_2 = 2\chi_{n-2,2} + 3\chi_{n-1,1} = p^2 + 2t - 3$  such that  $g(\varepsilon) = 1$ . For  $3 \leq k \leq n-3$ , let  $\sigma_k$  be a permutation such that  $p(\sigma_k) = k$ ,  $t(\sigma_k) = 0$ , and let  $\eta_k$  be a permutation such that  $p(\eta_k) = 0$  and  $t(\eta_k) = k$ . Then,*

$$\max\{g(\sigma_k), g(\eta_k)\} \geq \frac{3k-5}{n^2 - kn + k + 2n - 5}.$$

PROOF. We can write

$$g = \alpha_1 \frac{p-1}{n-1} + \alpha_2 \frac{p^2 + 2t - 3}{n^2 - 3}$$

for some  $\alpha_1$  and  $\alpha_2$  such that  $\alpha_1 + \alpha_2 = 1$ . Then,

$$\begin{aligned} g(\sigma_k) &= \alpha_1 \frac{k-1}{n-1} + \alpha_2 \frac{k^2-3}{n^2-3} \quad \text{and} \\ g(\eta_k) &= -\alpha_1 \frac{1}{n-1} + \alpha_2 \frac{2k-3}{n^2-3}. \end{aligned}$$

We observe  $g(\sigma_k)$  and  $g(\eta_k)$  are linear functions of  $\alpha_1$  and  $\alpha_2$  and that for

$$\alpha_1 = \frac{kn - k - 2n + 2}{-n^2 + kn - k - 2n + 5} \quad \text{and} \quad \alpha_2 = \frac{3 - n^2}{-n^2 + kn - k - 2n + 5}$$

we have

$$(10.4.1) \quad g(\sigma_k) = g(\eta_k) = \frac{3k-5}{n^2 - kn + k + 2n - 5} \quad \text{and} \quad \alpha_1 + \alpha_2 = 1.$$

Let

$$\lambda_1 = \frac{n^2 + 2kn - 3n - 2k}{k(3k-5)} \quad \text{and} \quad \lambda_2 = \frac{kn^2 - k^2n - n^2 + k^2 + 3n - 3k}{k(3k-5)}.$$

Then,  $\lambda_1, \lambda_2 > 0$  and

$$\lambda_1 g(\sigma_k) + \lambda_2 g(\eta_k) = \alpha_1 + \alpha_2 = 1.$$

Comparing this with (10.4.1), we conclude that there are no values  $\alpha_1, \alpha_2$  such that  $\alpha_1 + \alpha_2 = 1$  and

$$g(\sigma_k), g(\eta_k) < \frac{3k-5}{n^2 - kn + k + 2n - 5}. \quad \square$$

PROOF OF THEOREM 4.3. The proof is similar to that of Theorem 3.2 (see §9). Without loss of generality, we assume that the maximum value of  $f_0$  is 1 and is attained at the identity permutation  $\varepsilon$ . Let  $g$  be the central projection of  $f_0$ . We deduce that  $g$  is a linear combination of  $\chi_{n-1,1}$  and  $\chi_{n-2,2}$ , and hence a linear combination of  $g_1$  and  $g_2$ , and that  $g(\varepsilon) = 1$ . Let us choose  $3 \leq k \leq n-3$  and let  $X_k$  be the set of permutations  $\sigma$  such that  $p(\sigma) = k$  and  $t(\sigma) = 0$ , and let  $Y_k$  be the set of permutations  $\eta$  such that  $p(\eta) = 0$  and  $t(\eta) = k$ . As in the proof of Theorem 3.2, we have

$$|X_k| \geq \frac{1}{5} \frac{n!}{k!}.$$

To choose a permutation  $\eta \in Y_k$ , one has to choose  $k$  2-cycles in  $n!/((n-2k)!k!2^k)$  ways and a derangement without 2-cycles on the remaining  $n-2k$  points. Hence, by (6.6.2),

$$|Y_k| \geq \frac{1}{5} \frac{n!}{k!2^k}.$$

Let us choose a permutation  $\sigma \in X_k$  and a permutation  $\eta \in Y_k$  and let  $Z = X_k$  if  $g(\sigma) \geq g(\eta)$  and let  $Z = Y_k$  otherwise. Then,

$$|Z| \geq \frac{n!}{5k!2^k}$$

and by Lemma 10.4,

$$g(\sigma) \geq \frac{3k-5}{n^2 - kn + k + 2n - 5} \quad \text{for all } \sigma \in Z.$$

The proof now proceeds as in the proof of Theorem 3.2, (§9).  $\square$

**11. The general case: Proofs.** Let us choose a convenient basis in  $\text{span}\{\chi_{n-1,1}, \chi_{n-2,2}, \chi_{n-2,1,1}\}$ :

$$\begin{aligned} g_1 &= \chi_{n-1,1} = p-1, & g_2 &= \chi_{n-2,2} + \chi_{n-2,1,1} + 3\chi_{n-1,1} = p^2-2, & \text{and} \\ g_3 &= \chi_{n-2,1,1} - \chi_{n-2,2} = 1-2t. \end{aligned}$$

11.1. DEFINITION. Let  $K$  be the set of all functions  $g \in \text{span}\{g_1, g_2, g_3\}$  such that  $g(\varepsilon) \geq g(\sigma)$  for all  $\sigma \in S_n$ . We call  $K$  the *central cone*.

Identifying  $\text{span}\{g_1, g_2, g_3\}$  with a 3-dimensional vector space  $\mathbb{R}^3$ , we see that conditions  $g(\varepsilon) \geq g(\sigma)$  define the central cone  $K$  as a convex polyhedral cone in  $\mathbb{R}^3$ . The condition  $g(\varepsilon) = 1$  defines a plane  $H$  in  $\mathbb{R}^3$  and the intersection  $B = H \cap K$  is a *base* of  $K$ , that is, a polygon such that every  $g \in K$  can be uniquely represented in the form  $g = \lambda h$  for some  $h \in B$ .

Our goal is to determine the structure of  $K$ . This is somewhat more complicated than in the 2-dimensional situations of §§9 and 10.

11.2. LEMMA. *Let us define functions*

$$\begin{aligned} r_1 &= \frac{-np + n + p^2 - 2}{n-2}, \\ r_2 &= 1 - 2t, \\ r_3 &= \frac{2np - 3p - 2n - p^2 - 2t + 6}{n^2 - 5n + 6}, \\ r_4 &= \frac{p + 2t - 2}{n-2}, \quad \text{and} \\ r_{5o} &= \frac{-2np + 3p^2 - 3p + 2tn + n - 3}{n^2 - 2n - 3}. \end{aligned}$$



Then,

(1) If  $\varepsilon \in S_n$  is the identity, then

$$r_1(\varepsilon) = r_2(\varepsilon) = r_3(\varepsilon) = r_4(\varepsilon) = r_{5o}(\varepsilon) = 1.$$

(2) If  $n$  is even, then  $r_1, r_2, r_3,$  and  $r_4$  are the vertices (in consecutive order) of the planar quadrilateral  $B = \text{conv}\{r_1, r_2, r_3, r_4\}$  that is a base of the central cone  $K$ .

(3) If  $n$  is odd, then  $r_1, r_2, r_3, r_4,$  and  $r_{5o}$  are the vertices (in consecutive order) of the planar pentagon  $B = \text{conv}\{r_1, r_2, r_3, r_4, r_{5o}\}$  that is a base of the central cone  $K$ .

PROOF. A function  $g \in \text{span}\{g_1, g_2, g_3\}$  can be written as a linear combination  $g = \alpha_1 g_1 + \alpha_2 g_2 + \alpha_3 g_3$ . Then,  $g(\varepsilon) = \alpha_1(n-1) + \alpha_2(n^2-2) - \alpha_3$  and the conditions  $g(\varepsilon) \geq g(\sigma)$  are written as

$$\alpha_1(n-1) + \alpha_2(n^2-2) + \alpha_3 \geq \alpha_1(p(\sigma)-1) + \alpha_2(p^2(\sigma)-2) + \alpha_3(1-2t(\sigma)),$$

which for  $\sigma \neq \varepsilon$  are equivalent to

$$\alpha_1 + \alpha_2(n+p(\sigma)) + \alpha_3 \frac{2t(\sigma)}{n-p(\sigma)} \geq 0.$$

Applying Lemma 6.8, we see that for even  $n$ , the system is equivalent to

$$(11.2.1) \quad \begin{aligned} \alpha_1 + n\alpha_2 &\geq 0, \\ \alpha_1 + (2n-3)\alpha_2 &\geq 0, \\ \alpha_1 + (2n-2)\alpha_2 + \alpha_3 &\geq 0, \\ \alpha_1 + n\alpha_2 + \alpha_3 &\geq 0, \end{aligned}$$

whereas for odd  $n$ , the system is equivalent to

$$(11.2.2) \quad \begin{aligned} \alpha_1 + n\alpha_2 &\geq 0, \\ \alpha_1 + (2n-3)\alpha_2 &\geq 0, \\ \alpha_1 + (2n-2)\alpha_2 + \alpha_3 &\geq 0, \\ \alpha_1 + (n+1)\alpha_2 + \alpha_3 &\geq 0, \\ n\alpha_1 + n^2\alpha_2 + (n-3)\alpha_3 &\geq 0. \end{aligned}$$

The set of all feasible 3-tuples  $(\alpha_1, \alpha_2, \alpha_3)$  is a polyhedral cone, which, for even  $n$ , has at most 4 extreme rays and for odd  $n$  has at most 5 extreme rays. We call an inequality of (11.2.1)–(11.2.2) *active* on a particular tuple if it holds with equality.

It is readily verified that for even  $n$  the following tuples span the extreme rays of the set of solutions to (11.2.1):

$$\begin{aligned} (-n, 1, 0) & \text{ 4th and 1st inequalities are active,} \\ (0, 0, 1) & \text{ 1st and 2nd inequalities are active,} \\ (2n-3, -1, 1) & \text{ 2nd and 3rd inequalities are active,} \\ (1, 0, -1) & \text{ 3rd and 4th inequalities are active,} \end{aligned}$$

and that for odd  $n$  the following tuples span the extreme rays of the set of solutions to (11.2.1):

$$\begin{aligned} (-n, 1, 0) & \text{ 5th and 1st inequalities are active,} \\ (0, 0, 1) & \text{ 1st and 2nd inequalities are active,} \\ (2n-3, -1, 1) & \text{ 2nd and 3rd inequalities are active,} \\ (1, 0, -1) & \text{ 3rd and 4th inequalities are active,} \\ (-2n-3, 3, -n) & \text{ 4th and 5th inequalities are active.} \end{aligned}$$

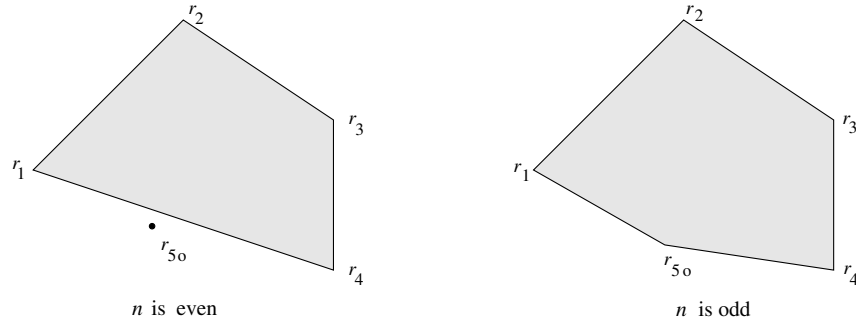


FIGURE 6. The base of the central cone.

We obtain  $r_1, r_2, r_3, r_4$ , and  $r_{5o}$  by scaling the corresponding linear combinations  $\alpha_1 g_1 + \alpha_2 g_2 + \alpha_3 g_3$  so that the value at the identity is equal to 1, and hence  $r_1, r_2, r_3, r_4$ , and  $r_{5o}$  lie on the same plane in  $\text{span}\{g_1, g_2, g_3\}$ .  $\square$

11.3. REMARK. We observe that  $r_1$  and  $r_{5o}$  have spike distributions (in particular,  $r_1$  has the “sharp spike” distribution of §5.2) corresponding to the cases of  $\gamma_1 = 1$  and  $\gamma_5 = 1$ , respectively, in Theorem 5.1; that  $r_2$  has the damped oscillator distribution corresponding to the case of  $\gamma_2 = 1$ ; and that  $r_3$  and  $r_4$  have bullseye distributions corresponding to the cases of  $\gamma_3 = 1$  and  $\gamma_4 = 1$ , respectively. If  $n$  is even, then  $r_{5o} \notin K$ , for if  $\sigma$  is a product of  $n/2$  commuting 2-cycles, so that  $p(\sigma) = 0$  and  $t(\sigma) = n/2$ , then  $r_{5o}(\sigma) = (n^2 + n - 3)/(n^2 - 2n - 3) > 1 = r_{5o}(\varepsilon)$ ; see Figure 6.

We observe that function  $r_3$  coincides with function  $r_1$  of Lemma 10.2 (the symmetric QAP), and that function  $r_2$  coincides with function  $r_1$  of Lemma 9.2 (the pure QAP).

PROOF OF THEOREM 5.1. The proof is completely similar to the proof of Theorem 3.1 (see §9). We use Lemma 11.2 instead of Lemma 9.2.  $\square$

To prove Theorem 5.2, we need one preliminary result.

11.4. LEMMA. Let  $g$  be a linear combination of  $g_1 = p - 1$ ,  $g_2 = p^2 - 2$ , and  $g_3 = 1 - 2t$  such that  $g(\varepsilon) = 1$ . For  $2 \leq k \leq n - 2$ , let  $\sigma_k$  be a permutation such that  $p(\sigma_k) = k$  and  $t(\sigma_k) = 0$ , let  $\eta$  be a permutation such that  $p(\eta) = 0$  and  $t(\eta) = 1$ , and let  $\theta$  be permutation such that  $p(\theta) = t(\theta) = 0$ . Then

$$\max\{g(\sigma_k), g(\eta), g(\theta)\} \geq \frac{k-2}{n^2 - kn + k - 2}.$$

PROOF. We can write

$$g = \alpha_1 \frac{p-1}{n-1} + \alpha_2 \frac{p^2-2}{n^2-2} + \alpha_3 (1-2t)$$

for some  $\alpha_1, \alpha_2$ , and  $\alpha_3$  such that  $\alpha_1 + \alpha_2 + \alpha_3 = 1$ . Then,

$$g(\sigma_k) = \alpha_1 \frac{k-1}{n-1} + \alpha_2 \frac{k^2-2}{n^2-2} + \alpha_3,$$

$$g(\eta) = -\frac{\alpha_1}{n-1} - \alpha_2 \frac{2}{n^2-2} - \alpha_3,$$

$$g(\theta) = -\frac{\alpha_1}{n-1} - \alpha_2 \frac{2}{n^2-2} + \alpha_3.$$

We observe that  $g(\sigma_k), g(\eta)$ , and  $g(\theta)$  are linear functions of  $\alpha_1, \alpha_2$ , and  $\alpha_3$  and that for

$$\alpha_1 = \frac{k(1-n)}{n^2 - nk + k - 2}, \quad \alpha_2 = \frac{n^2 - 2}{n^2 - nk + k - 2}, \quad \text{and} \quad \alpha_3 = 0$$

we have

$$(11.4.1) \quad g(\sigma_k) = g(\eta) = g(\theta) = \frac{k-2}{n^2 - nk + k - 2} \quad \text{and} \quad \alpha_1 + \alpha_2 + \alpha_3 = 1.$$

Let

$$\lambda_1 = \frac{n^2 - 2n}{k^2 - 2k}, \quad \lambda_2 = \frac{n^2 - nk}{2k - 4}, \quad \text{and} \quad \lambda_3 = \frac{n^2 - kn - 2n + 2k}{2k}.$$

Then,  $\lambda_1, \lambda_2, \lambda_3 > 0$  and

$$\lambda_1 g(\sigma_k) + \lambda_2 g(\eta) + \lambda_3 g(\theta) = \alpha_1 + \alpha_2 + \alpha_3 = 1.$$

Comparing this with (11.4.1), we conclude that there are no values  $\alpha_1, \alpha_2,$  and  $\alpha_3$  such that  $\alpha_1 + \alpha_2 + \alpha_3 = 1$  and

$$g(\sigma_k), g(\eta), g(\theta) < \frac{k-2}{n^2 - nk + k - 2}. \quad \square$$

**PROOF OF THEOREM 5.2.** The proof follows those of Theorem 3.2 (see §9) and Theorem 4.3 (see §10) with some modifications. Let  $X_k$  be the set of all permutations  $\sigma$  such that  $p(\sigma) = k$  and  $t(\sigma) = 0$ . As in the proof of Theorem 3.2, we have

$$|X_k| \geq \frac{1}{5} \frac{n!}{k!}.$$

Let  $Y$  be the set of all permutations  $\sigma$  such that  $p(\sigma) = 0$  and  $t(\sigma) = 1$ . To choose a permutation  $\sigma \in Y$ , one has to choose a 2-cycle in  $\binom{n}{2}$  ways and then an arbitrary derangement on the remaining  $(n-2)$  symbols without 2-cycles. Using (6.6.2), we estimate

$$|Y| \geq \frac{1}{5} \frac{n!}{2(n-2)!} (n-2)! = \frac{1}{10} n!.$$

Let us choose a permutation  $\sigma_k \in X_k$ , a permutation  $\eta \in Y$ , and a permutation  $\theta \in X_0$ . Let us choose  $Z$  to be one of  $X_k, X_0,$  and  $Y$ , depending where the maximum value of  $g(\sigma_k), g(\eta),$  or  $g(\theta)$  is attained. Hence,

$$|Z| \geq \frac{n!}{5k!}$$

and by Lemma 11.4,

$$g(\sigma) \geq \frac{k-2}{n^2 - kn + k - 2} \quad \text{for all } \sigma \in Z.$$

We proceed now as in the proof of Theorem 3.2, (see §9).  $\square$

To obtain Corollary 5.4, we choose  $\gamma = 1/2$  and  $k = O(\alpha)$  in Part (1), and  $\gamma = 1/2$  and  $k = O(n^{1-\epsilon})$  in Part (2).

**12. Concluding remarks and open questions.** The estimates of Theorems 2.3, 3.2, 4.3, and 5.3 for the number of near-optimal permutations can be used to bound the optimal value by a sample optimum in branch-and-bound algorithms. Those estimates are (nearly) best possible for the generalized problem (1.2). However, it is not clear whether they can be improved in the case of standard QAP (see 1.1) or how to improve them in interesting special cases. In particular, we ask the following:

- *Question:* Let  $f: S_n \rightarrow \mathbb{R}$  be the objective function in the TSP (cf. §§2 and 3), let  $\bar{f}$  be the average value of  $f$ , and let  $f_0 = f - \bar{f}$ . Let  $\tau$  be an optimal permutation, so that  $f_0(\tau) \geq f_0(\sigma)$  for all  $\sigma \in S_n$ . Is it true that for any fixed  $\gamma > 0$  there is a number  $\delta = \delta(\gamma) > 0$  such that the probability that a random permutation  $\sigma \in S_n$  satisfies the inequality  $f_0(\sigma) \geq (\gamma/n)f_0(\tau)$  is at least  $n^{-\delta}$  for all sufficiently large  $n$ ?

Barvinok (2002) showed that this is indeed the case for an arbitrary QAP (see 1.1) provided  $f_0$  is replaced by its absolute value  $|f_0|$ .

As well as finding improved approximation algorithms, it would be very interesting to find corresponding hardness results with an ultimate goal of proving sharp bounds. No hardness of approximation results with respect to the average are known.

Another question is whether approximations can be obtained *deterministically*. The random sampling algorithm in the “bullseye” case of §2 can be relatively easily derandomized. Whether the same is true for algorithms of §§3–5 is not clear at the moment.

Our methods can be applied to study the distribution of values in the assignment problems of higher order and their special cases, such as the Weighted Hypergraph Matching Problem.

**Acknowledgments.** This research was partially supported by NSF Grant DMS 9734138.

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A. Barvinok: Department of Mathematics, University of Michigan, 1858 East Hall, Ann Arbor, Michigan 48109-1109; e-mail: barvinok@umich.edu

T. Stephen: Institute of Mathematics and Its Applications (IMA), University of Minnesota, 400 Lind Hall, 207 Church Street S.E., Minneapolis, Minnesota 55455-0436; e-mail: tamon@ima.umn.edu